

Lay Down the Common Metrics

Evaluating PoW Consensus Protocols' Security

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SUBCHAINS PUBLISH OR PERISH
TORTOISE AND HARES BYZCOIN GOSHAWK
BAHACK'S IDEA BITCOIN-NG (AETERNITY, WAVES)

BITCOIN'S NAKAMOTO CONSENSUS

ETHEREUM POW DECOR+ (ROOTSTOCK)
GHOST-DAG SPECTRE CHAINWEB
FRUITCHAINS PHANTOM BOBTAIL
THE INCLUSIVE PROTOCOL GHOST
CONFLUX

?

 *bitcoin*

2

1

Bitcoin's Nakamoto Consensus

NC



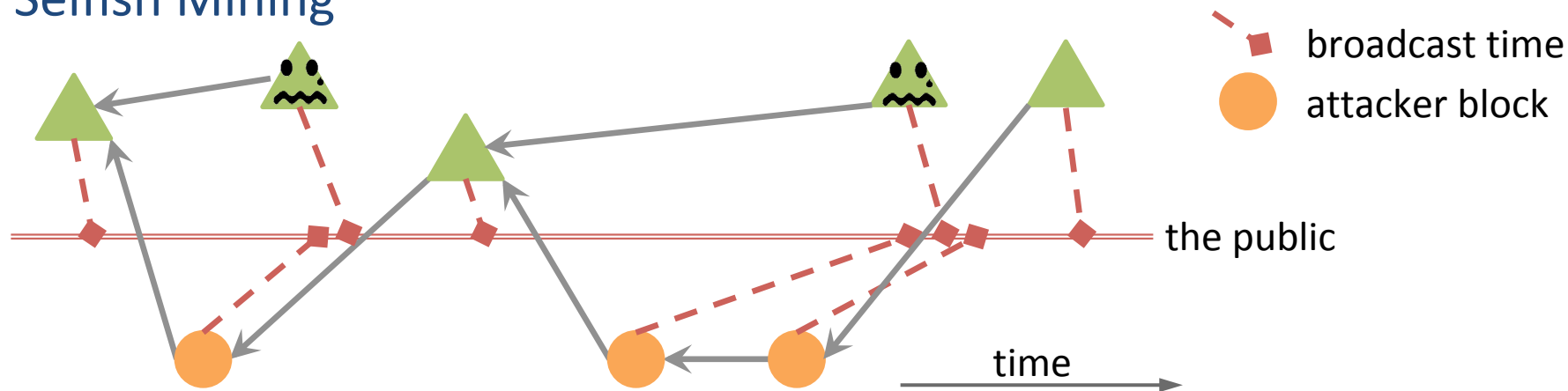
- To resolve fork
 - Longest chain (roughly) if there is one
 - First-received in a tie
- To issue rewards
 - Main chain blocks ▲ receive full rewards
 - Orphaned blocks △ receive nothing

Key Weakness

- Imperfect chain quality:
A <50% attacker can modify the blockchain with high success rate

Imperfect Chain Quality 🙌 3 Attacks

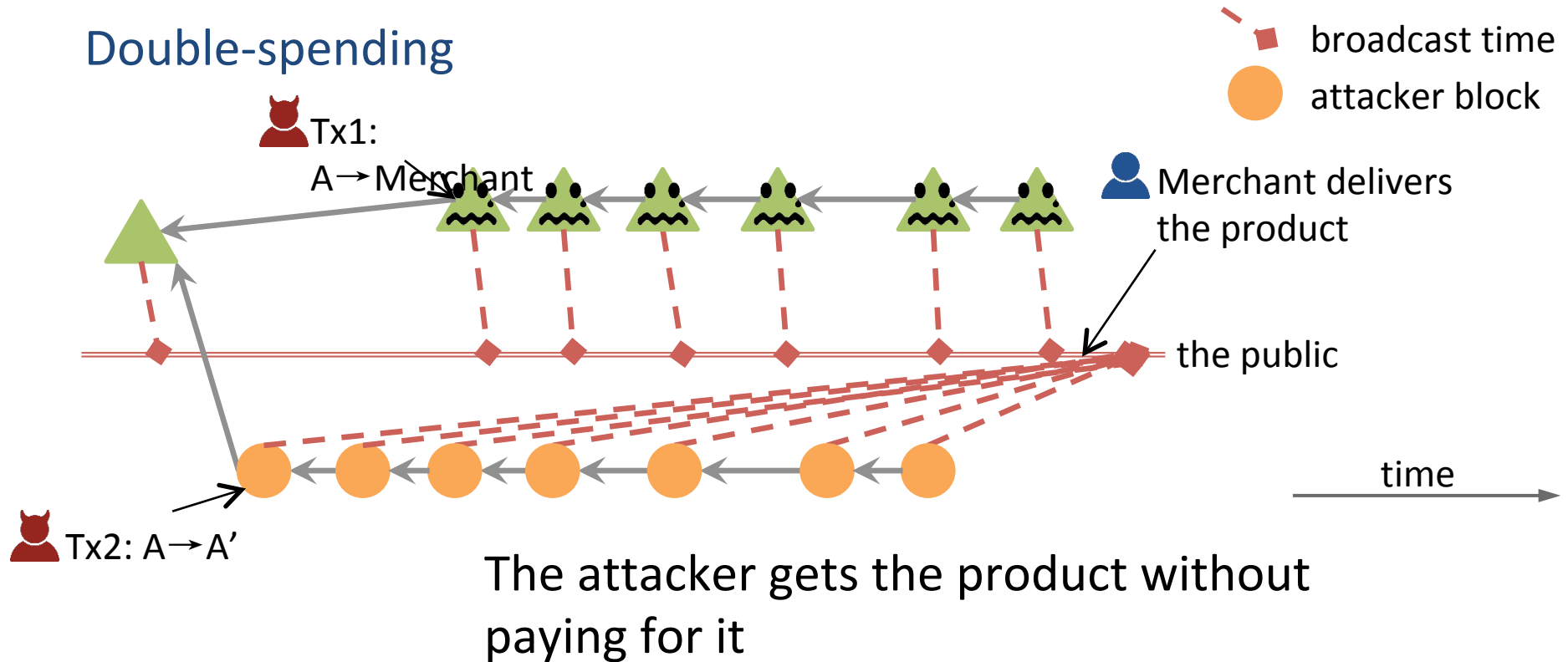
Selfish Mining



The attacker gains **unfair** block rewards; rational miners would join the attacker, which damages decentralization

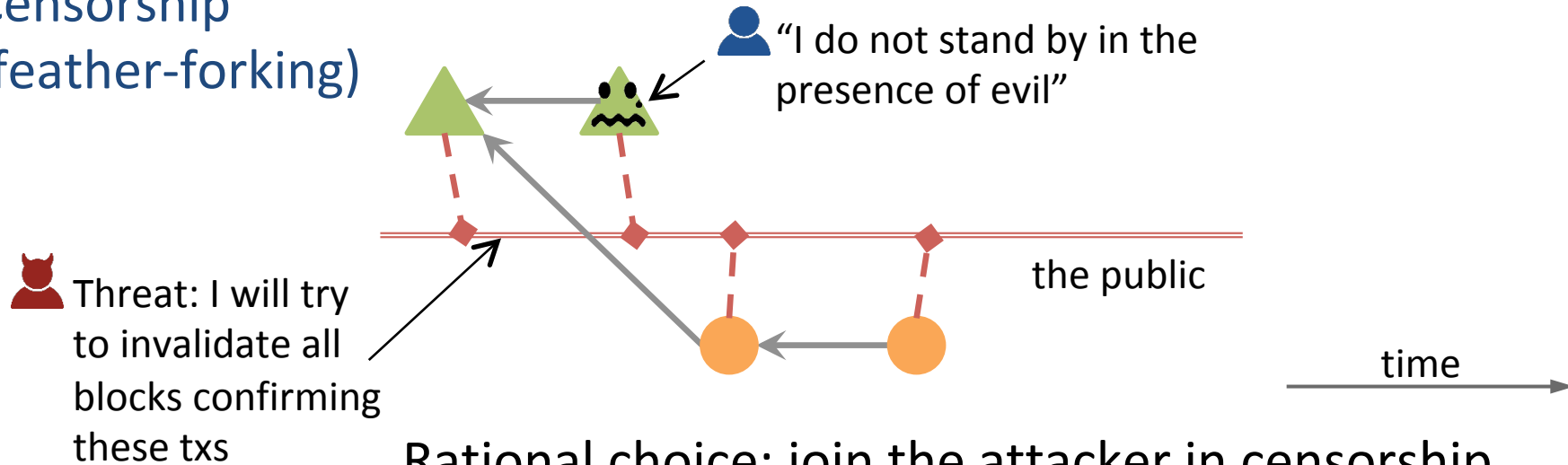
Imperfect Chain Quality 🙌 3 Attacks

Double-spending



Imperfect Chain Quality 🙅 3 Attacks

Censorship (feather-forking)



Rational choice: join the attacker in censorship
The attacker becomes a *de facto* owner

Our Evaluation Framework: 4 Metrics

A protocol claims to be more secure than NC:

it either ■ achieves better chain quality ① ②

or ■ resists better against all three attacks:

■ selfish mining 👉 incentive compatibility ①

■ double-spending 👉 subversion gain ①

■ censorship 👉 censorship susceptibility ②

(check the paper for the math definitions)

① profit-driven
adversary

② byzantine
adversary

Better-than-NC Candidates



Better-chain-quality protocols

“I can raise the chain quality”

- **UTB**: Ethereum PoW, Bitcoin-NG (Aeternity, Waves)
- **SHTB**: DECOR+ (Rootstock)
- **UDTB**: Byzcoin, Omniledger
- **Publish or Perish**

Attack-resistant protocols

“I don’t need to raise the chain quality, I can defend against the attacks”

- **Reward-all** (“compensate the losers”): **Fruitchains**, Ethereum PoW, Inclusive, SPECTRE, PHANTOM, ...
- **Punishment** (“fine all suspects”): **DECOR+**, Bahack’s idea
- **Reward-lucky** (content-based reward): **Subchains**, Bobtail

In this talk
Check the paper

MDP-based Method

Main idea

Model the protocol execution as a **Markov decision process (MDP)**, enumerate all the attacker's **reasonable** strategies, find the ones that optimize the metrics

Step 1

Define **the attacker's utility** according to the security metric of interest. e.g., in selfish mining:

$$\text{utility} = \text{attacker's rewards} / \text{all the rewards}$$

Step 2

Model the protocol as an MDP

MDP-based Method

- Step 3 Solve the MDP, compute the attacker's optimal strategies and their maximum utilities in various settings
- Step 4 Compare the utilities with NC, find out when they are better/worse
- Step 5 Check the respective strategies, find out why

Cows Are Not Round in Reality



Do not equate the security of a consensus protocol with its cryptocurrency

- Many **real-world factors** affect the attack difficulty (e.g., 51% attack against ETC vs. against Bitcoin)
- Several systems rely on **extra protection** for certain attack resistance



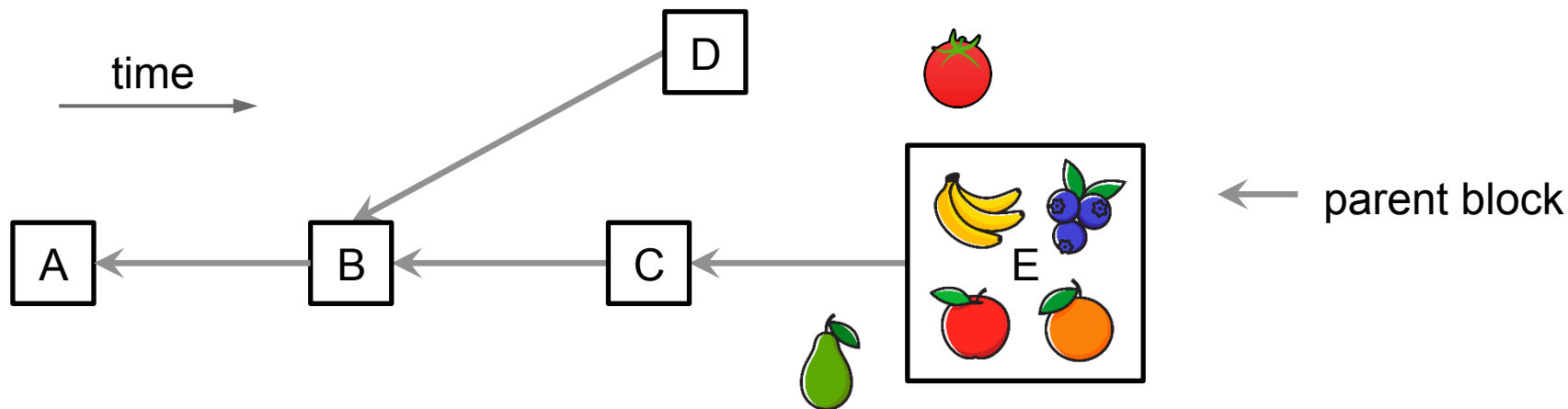
The Evaluation Results

Simplified Results

😊 better
😐 it depends
😞 worse

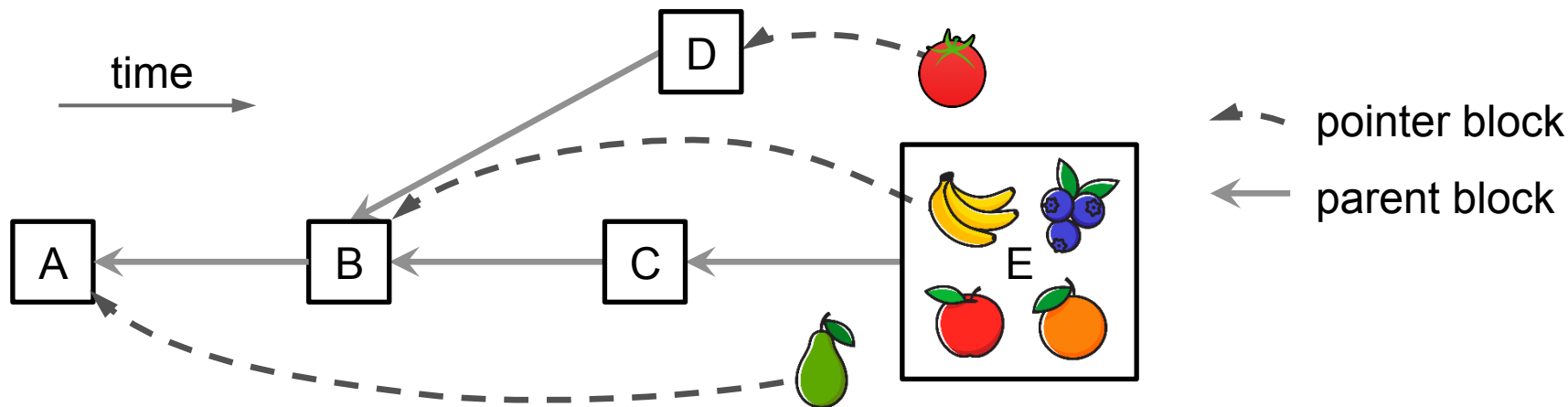
“Better-chain-quality”	Chain Quality	“Attack-resistant”	Incentive compatibility	Subversion gain	Censorship susceptibility
Uniform tie-breaking	😞	Reward-all 👉Fruitchains Punishment 👉Reward-splitting Reward-lucky 👉Subchains	😞	😞	😊
Smallest-hash tie-breaking	😞				
Unpredictable deterministic tie-breaking	😞		😊	😊	😞
Publish or perish	😞		😞	😞	😞

Attack-Resistant👉Reward-All: Fruitchains



- Same mining procedure, two products:
 - A **block** if the **first** k bits of $H(\text{candidate}) < D1$
 - A **fruit** if the **last** k bits of $H(\text{candidate}) < D2$
- Fruits in blocks; txs in fruits
- Fork-resolving: longest chain + first received (same as NC, RS and Subchains)

Attack-Resistant👉Reward-All: Fruitchains



- Each fruit has a **pointer block**: a recent block the fruit miner is sure will not be orphaned

A fruit is valid if

- The pointer block is in the main chain (sorry tomato)

And

- $\text{Gap}(\text{fruit}) = \text{height}(\text{host}) - \text{height}(\text{pointer}) < \text{TimeOut}$
(If $\text{TimeOut} = 3$, pear is hopeless)

Reward distribution

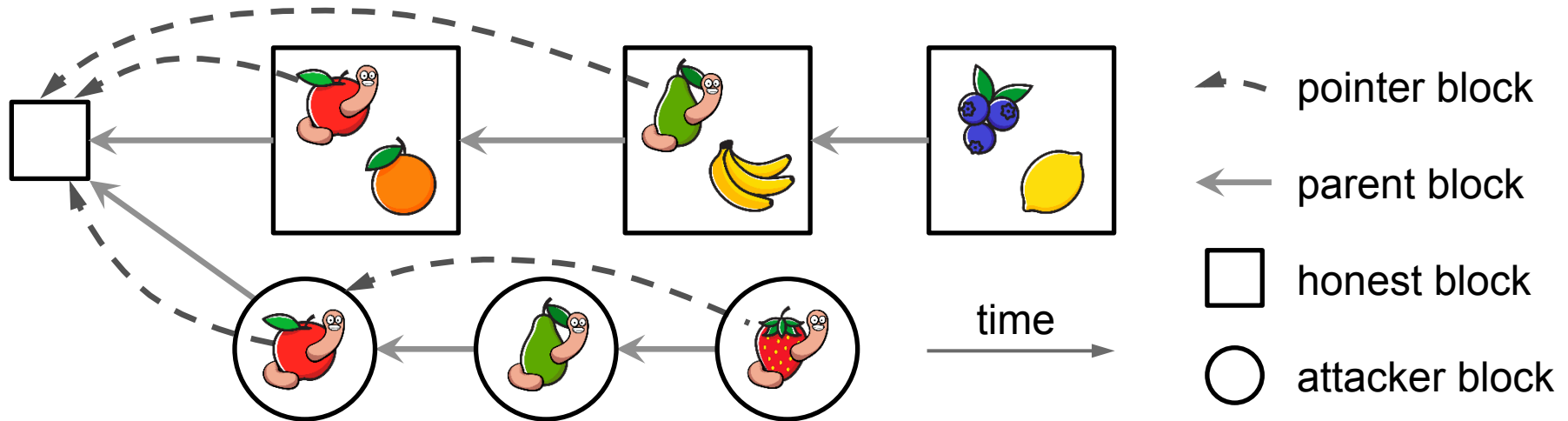
- Valid fruits receive rewards; blocks, nothing

Fruitchains Results

😊 better
😐 it depends
😞 worse

😞 Incentive
compatibility &
Subversion Gain

■ **Risk-free units** -> more audacious behaviors: attacker
uses **worthless blocks** to invalidate honest fruits;
attacker's first fruits are in both chains

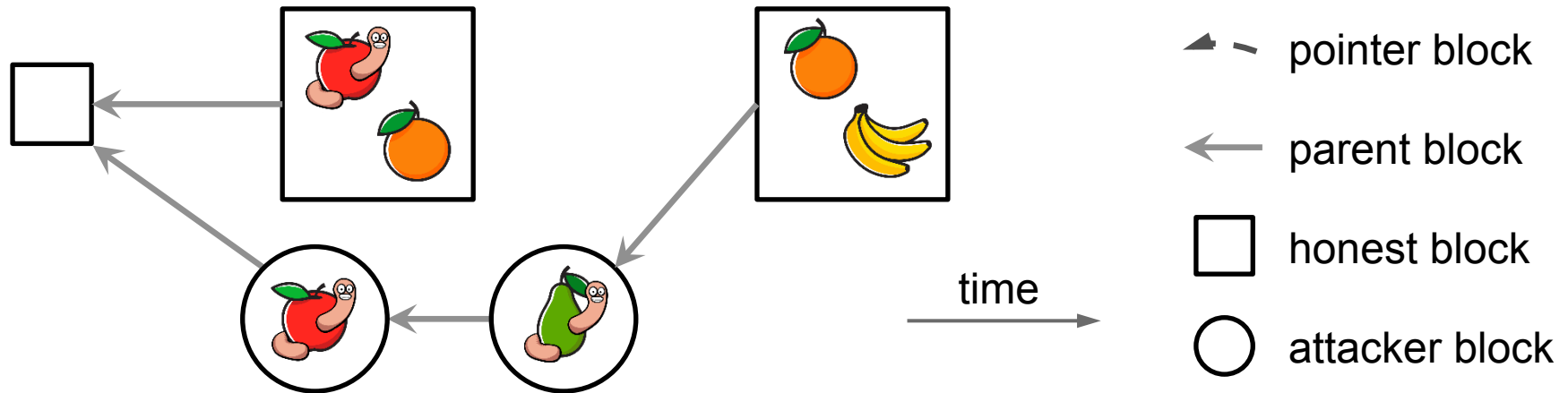


Fruitchains Results

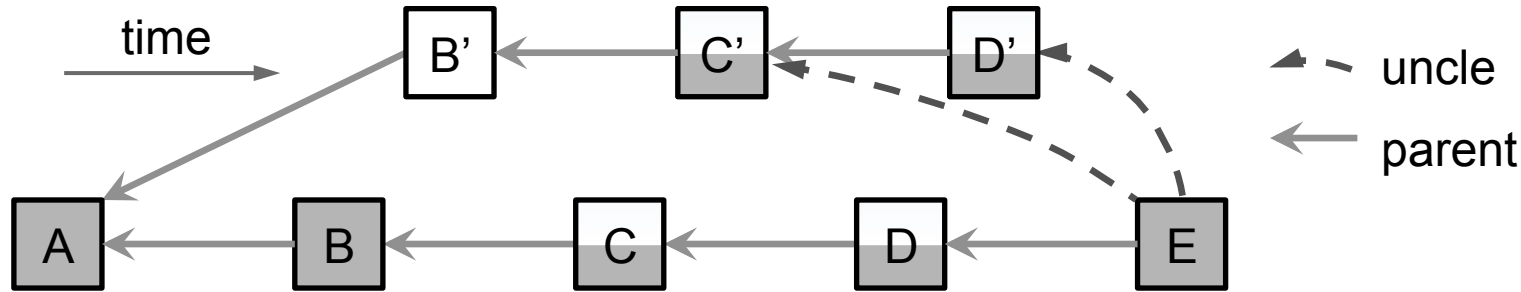
😊 better
😐 it depends
😞 worse

😊 Censorship
Susceptibility

- Fruits in invalidated blocks might be added back later (lucky orange)



Attack-Resistant👉Punishment: RS



- An **uncle** is valid if
 - $\text{Gap}(\text{uncle}) = \text{height}(\text{host}) - \text{height}(\text{uncle}) < \text{TimeOut}$
(B' is hopeless if TimeOut=3)
- Each block reward is **evenly split** among competing block & uncles of the same height

No pointer, unlike Fruitchains

(RS is modified from DECOR+, but their results are not the same!)

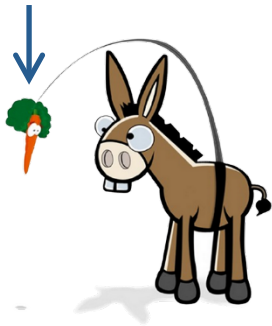
RS Results

😊 better
😐 it depends
😞 worse

😊 Incentive
compatibility &
Subversion Gain

- 3-confirmation RS performs better than 9-conf. Fruitchains

Subversion
Bounty



Min double-spending reward to incentivize double-spending attack attempts

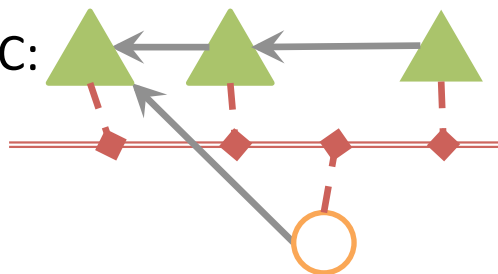
Attacker controls 10% mining power, 6-conf.,
bounty = 102 block rewards in NC,
 346 in RS,
 0 in Fruitchains

Censorship Susceptibility of RS

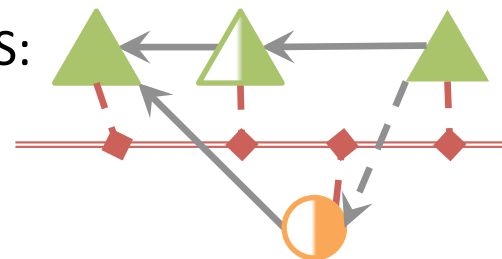


weak attackers

In NC:

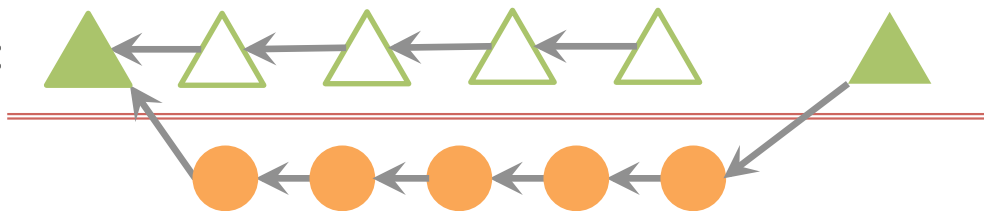


In RS:

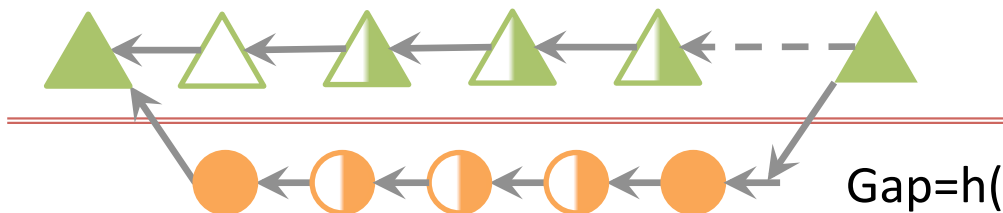


strong attackers

In NC:



In RS:



Gap= $h(\text{host}) - h(\text{self})$

Rewarding the Bad vs. Punishing the Good

A dilemma

When chain quality is not perfect ...

- Reward all -> no risk to double-spend
- Punish -> aid censorship
- Reward lucky -> lucky \neq good

Need to go beyond reward distribution policy to solve all attacks

Discussion

Simplicity is
beauty

- No protocol comprehensively outperforms NC

What not to do

- Designing protocols too complicated to analyze
- Security analysis
 - against one attack strategy
 - against one attacker incentive
 - with unrealistic parameters

Discussion

Better chain
quality & attack
resistance?

Practical assumptions

- Awareness of network conditions
- Loosely synchronized clock
- Real-world commitments

Outsource liability to raise attack resistance

- Introduce additional punishment rules (embed proofs of malicious behavior in blockchain)
- Solve at layer 2 (e.g. lightning guarantees double spending resistance)

Short Conclusion

- Tell anyone that claims to have a perfectly secure consensus protocol...



ACADEMIA IS WATCHING YOU

Thank you!

Code: github.com/nirenzang/PoWSecurity

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