SoK: Sanitizing for Security

Dokyung Song, Julian Lettner, Prabhu Rajasekaran, Yeoul Na, Stijn Volckaert, Per Larsen, Michael Franz

University of California, Irvine

Finding Bugs in C/C++



Finding Bugs in C/C++



Dynamic Analysis Tools for C/C++

More than 35 years of research in Dynamic Analysis Tools – often-called "Sanitizers"

 that find vulnerabilities specific to C/C++

Oscar

					Undangle	FreeSentry He	exType
					SoftBounds+CETS	SGXBour	nds TySan
					Dr. Memory	CaVer Effe	ctiveSan
	Purify			MSCC	LBC	TypeSa	n CUP
	Electric Fence	;		Memcheck	PAriCheck UBSa	an MSan Da	ngSan
Bcc	RTCC S	Safe-C P&F	PageHeap	CRED D&A	BBC ASan	DangNull Lov	w-Fat CRCount
1980	1990	1995	2000	2005	2010	2015	2019

Exploit Mitigation vs. Sanitization (1/2)

Attack Flow

Integer Overflow — *Heap Overflow* + *Function Pointer Overwrite* — *Indirect Call*





Exploit Mitigation vs. Sanitization (2/2)

	Exploit Mitigation	Sanitization
The goal is to	Mitigate attacks	Find vulnerabilities
Used in	Production	Pre-release
Performance budget is	Very limited	Much higher
Policy violation leads to	Program termination	Problem diagnosis
Violations triggered at location of bug	Sometimes	Always
Tolerance for FPs is	Zero	Somewhat higher
Surviving benign errors is	Desired	Not desired

Undefined Behavior in C/C++

. . .

- Buffer overflow —
- Use-after-free -
- Type errors
- Format string bug
- Signed integer overflow
- Null pointer dereferences

• etc.

J.2 Undefined behavior

The behavior is undefined in the following circumstances: ...

— Addition or subtraction of a pointer into, or just beyond, an array object and an integer type produces a result that does not point into, or just beyond, the same array object (6.5.6).

– An object is referred to outside of its lifetime (6.2.4).

— A pointer is used to call a function whose type is not compatible with the referenced type.

— An object has its stored value accessed other than by an lvalue of an allowable type (6.5).

Undefined Behavior in C/C++



. . .

Security Implications of Undefined Behavior in C/C++ (1/2)

1. Memory and type safety violations vulnerable to memory exploits



Security Implications of Undefined Behavior in C/C++ (2/2)

2. Compilation of a program having UBs may result in vulnerable code



Security Implications of Undefined Behavior in C/C++ (2/2)

2. Compilation of a program having UBs may result in vulnerable code



Security Implications of Undefined Behavior in C/C++ (2/2)

2. Compilation of a program having UBs may result in vulnerable code



Low-Level Vulnerabilities in C/C++ (1/2)



• **Most of these vulnerabilities** can manifest as memory and type safety violations.

Low-Level Vulnerabilities in C/C++ (2/2)



- **Some UBs** may lead to unsafe code generation today.
- And, *things can change* as compiler optimizations evolve called timebombs*.

^{*} W. Dietz, P. Li, J. Regher, and V. Adve; "Understanding integer overflow in C/C++." In *ICSE*, 2012.

Low-Level Vulnerabilities in C/C++ (2/2)



- **Some UBs** may lead to unsafe code generation today.
- And, *things can change* as compiler optimizations evolve called timebombs*.

^{*} W. Dietz, P. Li, J. Regher, and V. Adve; "Understanding integer overflow in C/C++." In *ICSE*, 2012.

Sanitizer Design and Implementation



Bug Finding Technique



Sanitizer Design and Implementation: Bug Finding Techniques

Bug Finding Techniques



Sanitizer Design and Implementation: Program Instrumentation

Program Instrumentation

Inlined Reference Monitor: Fine-grained run-time monitoring of program behavior to detect bugs as they occur. Intermediate Representation Source code Binary External (C/C++)(e.g., LLVM IR) call⁄ Libraries 11101110101 call check check(); 01010101101 (d) Library Interposition 01010101011 *ptr = 3;store ptr 0101 Compiler Compiler LD PRELOAD for instrumenting only Frontend Backend calls to dynamically-linked external library functions (b) IR-level (a) Language-level (c) Binary-level

Sanitizer Design and Implementation: Metadata Management



May 2019

Sanitizer Design and Implementation: Precision and Overheads



Our Analysis of Sanitizers

- Our analysis of 37 tools
 - We benchmarked 10 publicly available sanitizers on the same experimental platform (<u>https://github.com/securesystemslab/</u>

sanitizing-for-security-benchmarks)

- Main observations
 - Performance is not a primary concern
 - Many false positives (marked as) in tools other than widely-used ones such as ASan
 - Most () ly have partial coverage of bugs ()
 - Widely deployed tools such as ASan have even smaller coverage

Sanitizers		IV. Bug Finding Techniques								V. Instr.				VI. Metadata Mgmt.							VIII. Analysis												
	Year Published Actively Maintained	Red-zone Insertion	Guard Pages	Reuse Delay	Per-pointer Bounds Tracking	Per-object Bounds Tracking	Lock-and-key Dangling Pointer Tagging	Uninit. Mem. Read Detection	Uninit. Value Use Detection	Pointer Casting Monitor	Pointer Use Monitor	Variadic Arg. Mismatch Detection	Stateless Monitoring	Language-level Instr.	[R-level Instr.	Binary Instr.	Library Interposition	Embedded Metadata	Direct-mapped Shadow	Multi-level Shadow	Custom Data Structure	Fat/Tagged Pointers	Disjoint Per-pointer Metadata	Statte Pretauata Suntial Safatti Vialation (III A 1)		temporal Safety Violation (III-A2) Use of Uninit. Variables (III-B)	Bad-casting (III-C)	Load Type Mismatch (III-C)	Func. Call Type Mismatch (III-C)	Variadic Func. Misuse (III-D)	Signed Integer Overflow (III-E)	Performance Overhead	Memory Overhead
Purify [27]	·92	V		~	<u>L</u>	<u>д</u> , ,			<u> </u>	<u> </u>	д,	-	S	-	-	E V	Г	щ	~	4	0		- 0	2 0			ш	-	щ	-	SC		
Memcheck [28]	'05 V			v					V							2				~												Ö	
Dr. Memory [29]	'11 V			v					v							~				~						DO						0	
LBC [30]	11	2												~						2					D	UU							
ASan [31]	12 12 V	-		~											1				1						-	D						C	*
Electric Fence [32]	'93		~	v													~		•						-	Ď						n/a	
PageHeap [33]	'00		-	v													~															n/a	
D&A Dangling [35]	'06		~												~		•																
Oscar [36]	'17		~														V	V															
RTCC [45]	'92				~												•	•															
Safe-C [46]	'94				~		~							~								~	/										n
P&F [47]	'97				1		~							1											-								
MSCC [52]	'04				~		~							~																			
SoftBound+CETS [48], [56]	'10				~									~	~										4								
Intel Pointer Checker [49]	'12 v	/			~										~																		
SGXBounds [54]	12 V 17				v										v			~								0							
CUP [55]	17			V	~										~			•				~	/			0							n
J&K [34]	'97				•	~															~) r
CRED [37]	'04					~								~							~) n
D&A Bounds [38]	'06					v																											r
BBC [39]	'09					~									~						~												
PAriCheck [41]	'10					v																~											r
Low-fat Pointer [42], [43]	'17 v	/				~								~	1						~	1											*
Undangle $[57]$	17					•										~						•			Υ.								
FreeSentry [59]	12							,						./								'											r
DangNull [58]	15 15													~						~	~					-							
DangSan [60]	13						~	,							v																		*
MSan [66]	17 15 v														v					~	r												*
CaVer [69]	15 V								~	1				1	v												-						
TypeSan [70]	15 16									1				~	v					~	*												*
HexType [71]	10									4				~	~					*	~												*
Loginov et al. [72]	'01										~			~					~								-					0	
LlVM TySan [73]	'17 v	/									~			~	1				~													0	1
EffectiveSan [74]	17 V 18										2			~	v			v				~						-				n/a	
Clang CFI [68]	10 15 V	/								1	~			~	v			V				*				U							*
HexVASAN [79]	15 V 17										~				~						~								•				*
ACAVASAIN 1/9	'12 v											~			~			V			~								-	•	-		





Programs conforming to the ISO standard









K. Memarian, J. Matthiesen, J. Lingard, K. Nienhuis, D. Chisnall, R. N. M. Watson, and P. Sewell. Into the Depths of C: Elaborating the De Facto Standards. In PLDI'16



Reducing Precision Gaps (1/2): Standard Compatibility

Compatibility with the ISO and de facto standards









Programs conforming to the de facto standard



Reducing Precision Gaps (1/2): Standard Compatibility

- Many programs transiently construct OOB pointers (de facto standards)
 - Thus, supporting this code idiom increases a tool's applicability
- Many tools, however, do not permit transient construction of OOB pointers
 - Some bounds checking tools invalidate pointers as soon as they go OOB
 - Dangling pointer tagging tools may incorrectly invalidate pointers, if OOB pointers are transiently stored in memory
- Many tools only support ISO standard compatibility by adding one byte between objects this is not enough in practice

Reducing Precision Gaps (2/2): Finding Elusive Bugs

Finding bugs that elude existing or widely-deployed sanitizers



Reducing Precision Gaps (2/2): Finding Elusive Bugs

- Subclasses of memory safety violations that elude AddressSanitizer:
 - Intra-object buffer overflow
 - Buffer overflow into a valid but unintended object
 - Uses of freed memory that are being reused

- Type errors beyond bad casting (static_casts)
 - C programs or C++ programs using C-style casts and reinterpret_casts
 - Type errors are UBs that may silently break programs if does not instruct the compiler using flags like -fno-strict-aliasing

Reducing Precision Gaps (2/2): Finding Elusive Bugs

- Finding these elusive bugs, in general, requires more precise dynamic metadata tracking
 - Tracking **per-pointer metadata** such as pointer bounds
 - Tracking **effective types** of object storage in memory
- However, such metadata tracking poses precision and performance challenges
 - C's weak type safety (e.g., pointer to integer casts, uses of void pointers) makes pointer metadata tracking difficult
 - The C standard has complex effective type and aliasing rules
- More research is needed in developing sanitizers that can find these elusive bugs while maintaining good compatibility

Pointer Metadata Tracking Challenges: Uninstrumented Code



Fat pointer is not compatible with uninstrumented code

Disjoint pointer metadata can get outdated, when uninstrumented code updates a pointer without updating corresponding metadata

Pointer Metadata Tracking Challenges: Pointer to Integer Casts

• Even with full instrumentation (via, e.g., dynamic binary translation), sanitizers can break programs having **pointer to integer casts**

some_object_type $* \rightarrow$ uint64_t

- Incompatible with fat/tagged pointers
- Disjoint pointer metadata can be a choice, but full pointer flow tracking across casts between pointers and integers can be expensive
- Existing tools stop tracking pointer metadata once they are cast to integers

Pointer Metadata Tracking Challenges: Multi-threaded Programs

- *Race-free* programs that use atomic operations can be problematic
 - Concurrent atomic operations from different threads without putting corresponding metadata operations into the same atomic unit can make metadata go out-of-sync
- Example of **naïve instrumentation**:



Type Error Detection Challenges

- Rules about determining an effective type of a stored value (i.e., **effective type rules**) are complex, due to weakly-typed nature of C
 - Prevalent uses of void pointer type and (omnipotent) character pointer type
 - malloc returns void *
 - memcpy-family of functions take and return void *
 - Type punning through C-style casts and union
- There are some tools that implement over-approximations of effective type rules, but precision and performance trade-offs are yet to be explored.
- Also, type error checking itself can be costly, because C's **aliasing rules** permit a stored value to be accessed by using pointers of many different types

Other Future Research Directions

Composing sanitizers

• Can find bugs closer to their source without generating duplicated bug reports for the bug's side-effects

Using hardware features to improve performance and compatibility

• e.g., Pointer tagging/memory tagging support in HW

Kernel and bare metal support

• Sanitizers for OS kernels, or non-user-space programs in general

Q & A

Thank you!

Dokyung Song Ph.D. Student at UC Irvine dokyungs@uci.edu