Keeping Authorities "Honest or Bust" with Decentralized Witness Cosigning

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IEEE Security & Privacy - May 24, 2016

We depend on many authorities

Conceptually simple but security-critical services

Time Services (NTP)



Digital Notaries





Naming Authorites





Certificate Authorities





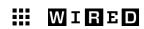
Randomness Authorities (e.g., Lotteries)



Software Update Services







Hack Obtains 9 Bogus Certificates for Prominent ...

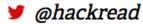
HACK OBTAINS 9 BOGUS CERTIFICATES FOR PROMINENT WEBSITES; TRACED TO IRAN



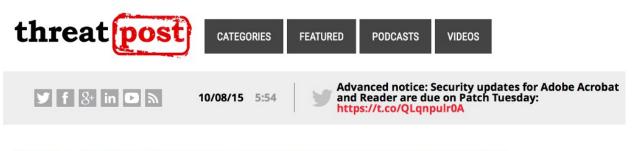
CYBER CRIME | SCAMS AND FRAUD

This Dude Hacked Lottery Computers To Win \$14.3M Jackpot In U.S.









Welcome > Blog Home > Cryptography > D-Link Accidentally Leaks Private Code-Signing Keys



New attacks on Network Time Protocol can defeat HTTPS and create chaos

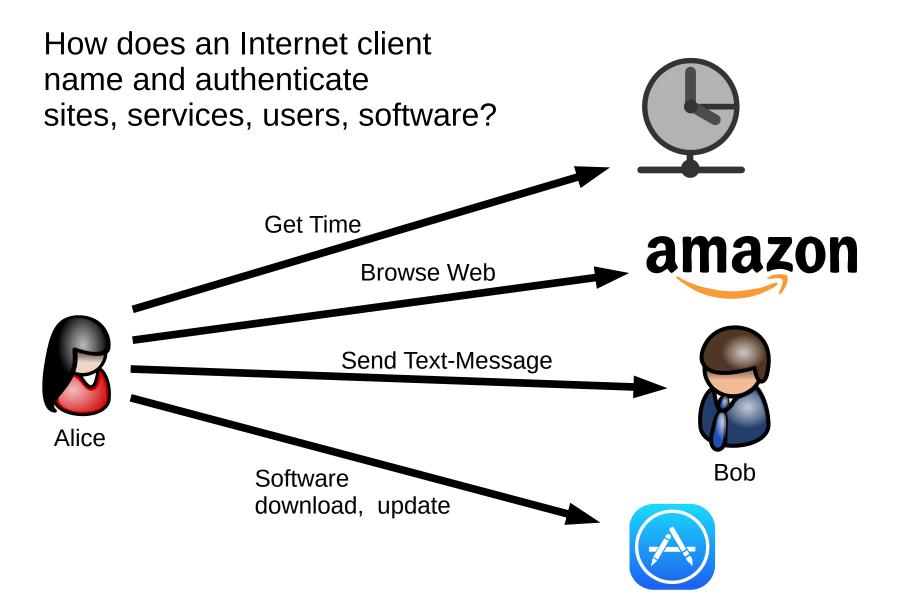
Exploits can be used to snoop on encrypted traffic and cause debilitating outages.



Talk Outline

- The trouble with trusting authorities
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Deep Dependence on Authorities



Deep Dependence on Authorities



Respect my Authoritah!











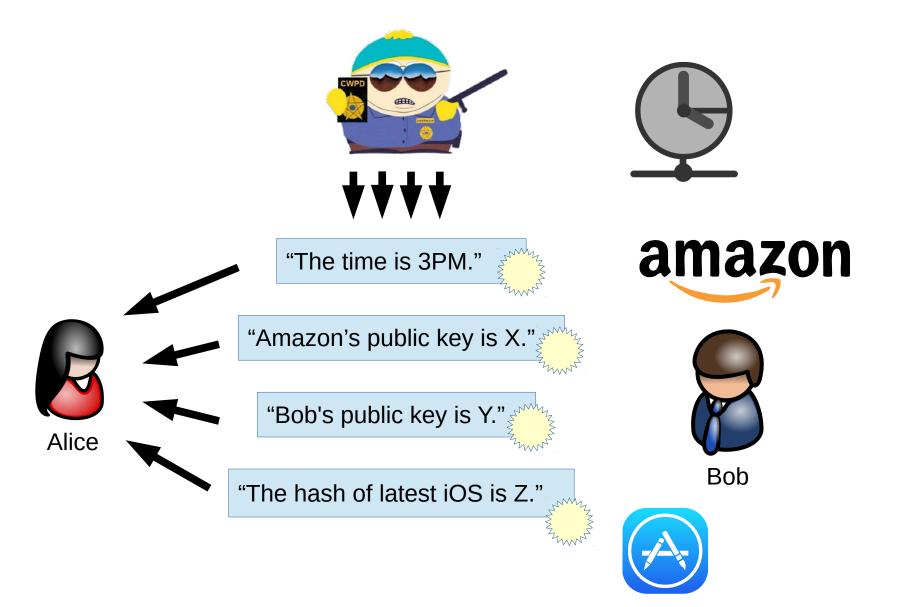


Alice

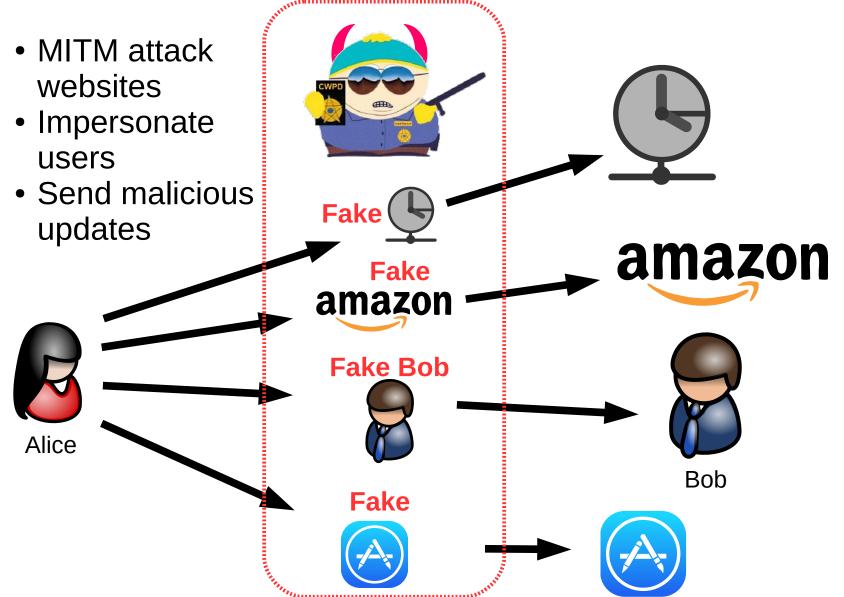
What is:

- The current time?
- Amazon's SSL public key?
- Bob's IM public key?
- Latest version of App?

Authorities Make & Sign Statements



Problem #1: Authority Compromise





Any CoA Can is issued to be a construction of the construction of



Attacker often needs to compromise only one

- wweakest-ink security
- @#\$% happens
 - DigiNotar,

 Oracle Commodo, 100 commodo, 100

CNNIC/MCS



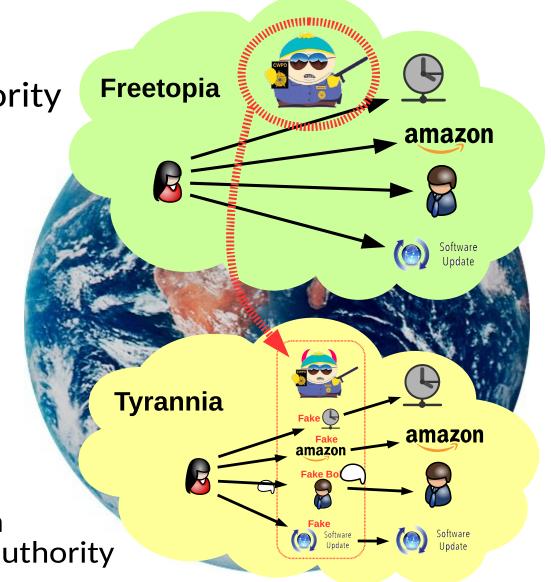
Problem #3: Secret Key Portability

 Attacker need not compromise authority "in-place"

 Instead steal the authority's secret key

- Use it to create an "evil twin" on attacker's turf
- Avoid detection
 by confining use
 to specific targets

 Block targets from reporting to real authority



Problem #4: Everybody Wants In

Hackers, organized crime, governments...



Security

Is Kazakhstan about to man-in-the-middle diddle all of its internet traffic with dodgy root certs?

Come on, guys. Don't go giving the Russians any ideas



Problem #4: Everybody Wants In

Hackers, organized crime, governments...



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What To Do?

We have to assume that no individual...

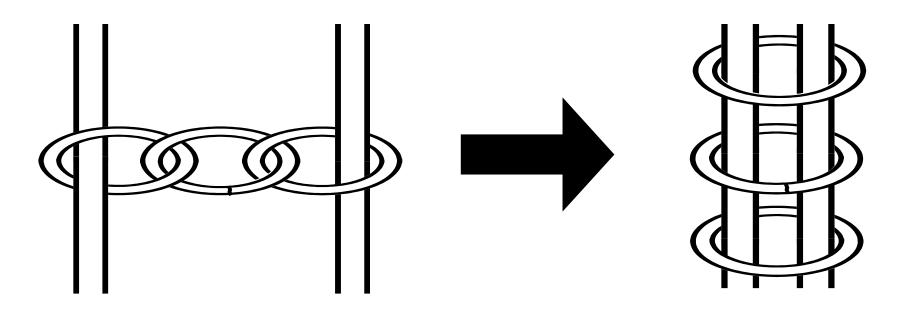
- Hardware platform
- Software system
- System/network administrator
- Authoritative organization

...is invulnerable to compromise (or coercion)

Decentralize the Authorities!

Split trust across independent parties

- So system resists compromise by individuals
- From weakest-link to strongest-link security
- Decentralized resistance to failure, coercion



Example: Tor Directory Authority

Split across ~10 servers – but is this enough?

Could attacker hack or coerce ~5 operators?

DIRECTORY AUTHORITIES

MORIA1 - 128.31.0.39 - RELAY AUTHORITY

TOR26 - 86.59.21.38 - RELAY AUTHORITY

DIZUM - 194.109.206.212 - RELAY AUTHORITY

TONGA - 82.94.251.203 - BRIDGE AUTHORITY

GABELMOO - 131.188.40.189 - RELAY AUTHORITY

DANNENBERG - 193.23.244.244 - RELAY AUTHORITY

URimage Credit: Jordan Wright) ORITY

MAATUSKA - 171.25.193.9 - RELAY AUTHORITY

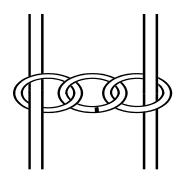
FARAVAHAR - 154.35.175.225 - RELAY AUTHORITY

LONGCLAW - 199.254.238.52 - RELAY AUTHORITY

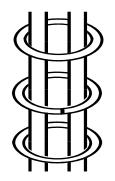
Trust-splitting needs to scale



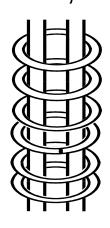
Weakest-link: T = 1



Strongest-link: T = 2-10



Collective authorities: T = 100s,1000s



Trust-splitting needs to scale

Must incorporate all diversity that makes sense

Not just ~10 parties "picked by someone"

Could we decentralize...

- TLS certificate validation and signing across the hundreds of certificate authorities?
- DNSSEC root zone maintenance and signing across the 1000+ worldwide TLD operators?
- A national cryptocurrency across the thousands of US national banks?

Make overall security **grow** as scale increase?

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Not Gonna Happen Overnight...



A First Step: **Transparency**

More readily achievable near-term

Who watches the watchers?
 Public witnesses!

Ensure that **any** authoritative statement:

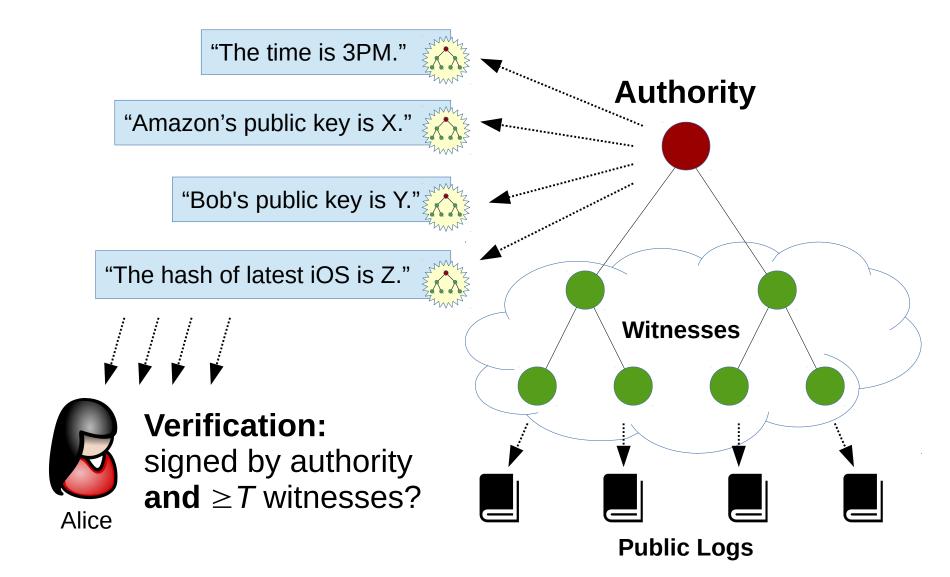
- Is exposed to public scrutiny
- Conforms to checkable standards

before clients will accept statement



Key: practical to "retrofit" existing authorities

Decentralized Witness Cosigning



Is the Signed Statement "Good"?

In general, witnesses don't (and can't) know for sure

- Does public key X really belong to Bob?
- Does software image Y have a secret backdoor?

But witnesses can still ensure all signatures are public

- If authority coerced or its keys used to produce bad statement, attacker can't ensure its secrecy
 - Backdoors possible but must "hide in plain sight"
- Creates "Ulysses Pact" deterrent against coercion
 - "the point...is to keep governments from even trying to put secret pressure on tech companies, because the system is set up so that the secret immediately gets out"
 Cory Doctorow, 10-March-2016

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Setup: Keypairs and CoSi Groups

Individual Keypairs:

Standard Schnorr (Ed25519)

- Private key: k
- Public key: K = g^k

CoSi group:

List of public keys

• K₁, K₂, ..., K_N

Assumptions:

- Verifier has full list
 - (nonessential)
- All keys self-signed
 - (important to avoid related-key attacks)

Schnorr Signature

- Generator g of prime order q group
- Public/private key pair: (K=g^k, k)

	Signer		Verifier
Commitment	V=g ^v	>	V
Challenge	С		c = H(M V)
Response	r = (v - kc)	>	r

Signature on M: (c, r)

Commitment recovery
$$V' = g^r K^c = g^{v-kc} g^{kc} = g^v = V$$
 Challenge recovery
$$c' = H(M|V')$$
 Decision
$$c' = c ?$$

Schnorr Multisignature

• Key pairs: $(K_1=g^{k_1}, k_1)$ and $(K_2=g^{k_2}, k_2)$

	Signer 1	Signer 2	veriner	
Commitment	$V_1=g^{V_1}$	$V_2 = g^{v_2}$	V_{1}	$V_2 V=V_1^*V_2$
Challenge	c c	<	$c = H(M V_1)$	c = H(M V)
Response	$r_1 = (v_1 - k_1 c$	$-\frac{r_2}{r_2} = (v_2 - v_2) k_2 c$	c) r ₁	$r_2 r=r_1+r_2$

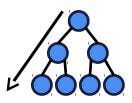
Signature on M: (c, r)) Same signature!

\/arifiar

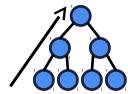
Commitment recovery Same verification! $V' = g^r K^c$ $K = K_1^* K_2$ Challenge recovery Done once! c' = H(M|V') Decision c' = c?

CoSi Protocol Signing Rounds

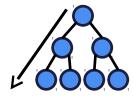
1. Announcement Phase



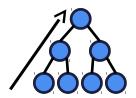
2. Commitment Phase



3. Challenge Phase



4. Response Phase

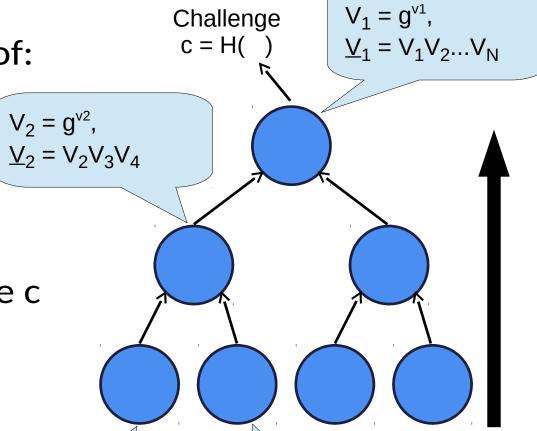


CoSi Commit Phase

Tree computation of:

- Commits V_i
- Aggregate commits <u>V</u>_i

Collective challenge c is hash of aggregate commit



$$V_3 = g^{v3},$$

$$\underline{V}_3 = V_3$$

$$V_4 = g^{V4},$$

$$\underline{V_4} = V_4$$

CoSi Response Phase

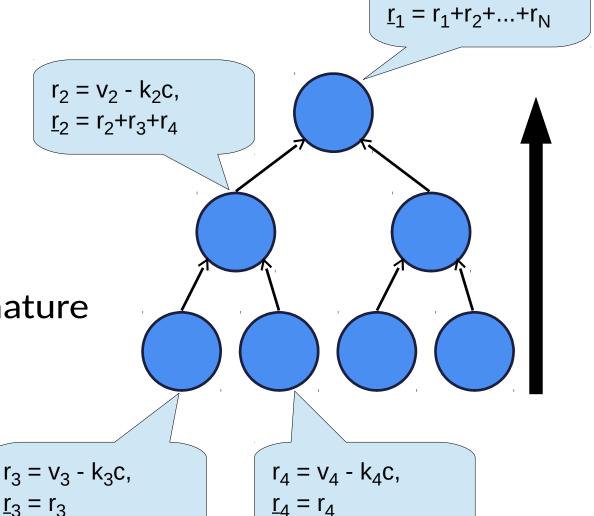
Compute

- Responses r_i
- Aggregate responses <u>r</u>;

Each (c,r_i) forms valid **partial** signature

(c,<u>r</u>₁) forms **complete**

signature



 $r_1 = v_1 - k_1 c$

Unavailable Witness Servers

Assume server failures are rare but non-negligible

Persistently bad servers get administratively booted

Exceptions: If a server A is down, proceed anyway

- Modified collective key: K'= K * K-1A
- Modified commitment: V'= V * V-1A
- Modified response: r'= r r_A

Verification: CoSi signature includes roll-call bit-vector

- Enables verifier to recompute modified public key K'
- Can use any criteria to decide if "too many" missing

Variations (see paper for details)

- Complex/contextual verification predicates
 - Witness subgroups, weights, expressions, ...
- Minimizing cothority certificate size
 - Via Merkle key-trees
- Tolerating network churn
 - Via binomial swap forests (Cappos, San Fermin)
- Tolerating cosigner churn
 - Avoiding restarts via commit trees
- Single-pass CoSi for asynchronous networks
 - Via BLS signatures, opportunistic signature combining

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Experimental Evaluation

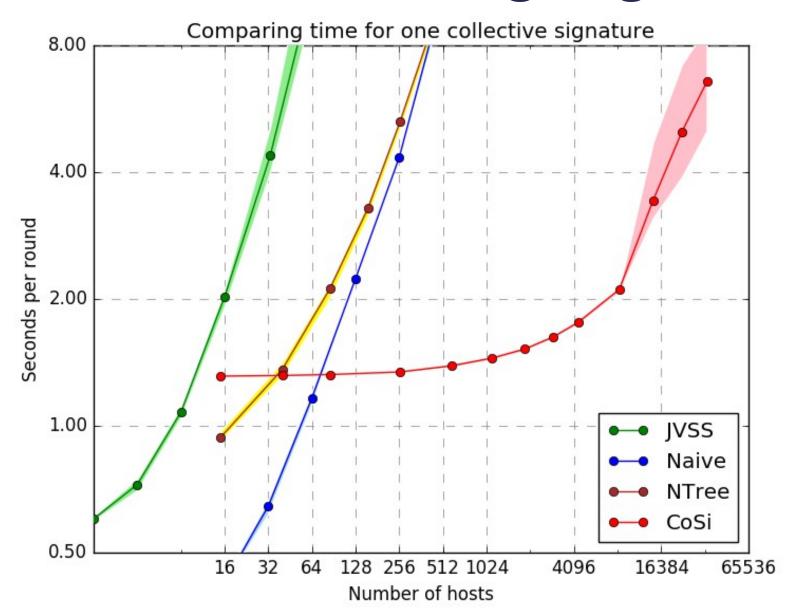
Experiments run on DeterLab network testbed

- Up to **32,768** virtual CoSi witnesses
- Multiplexed atop up to 64 physical machines
 - introduces oversubscription overhead, unfortunately
 - Conservative results, likely worse than "real" deployment
- Impose 200ms roundtrip latencies between all servers
 - to simulate globally-distributed witness group

Future: deploy, evaluate at scale on "real Internet"

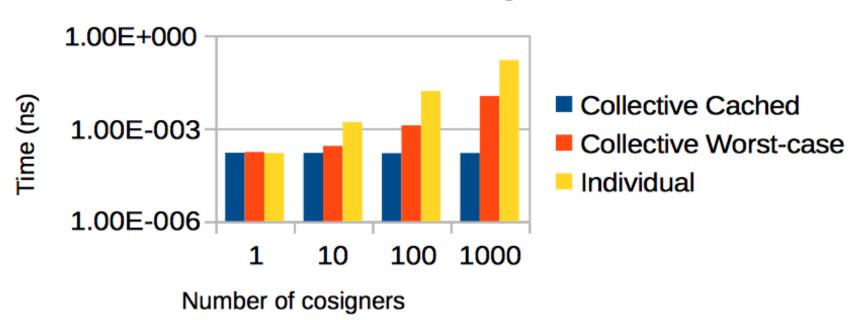
- Evaluate impact of high node, network churn
- See paper for approaches to handling if/when needed

Results: Collective Signing Time



Results: Verification Cost

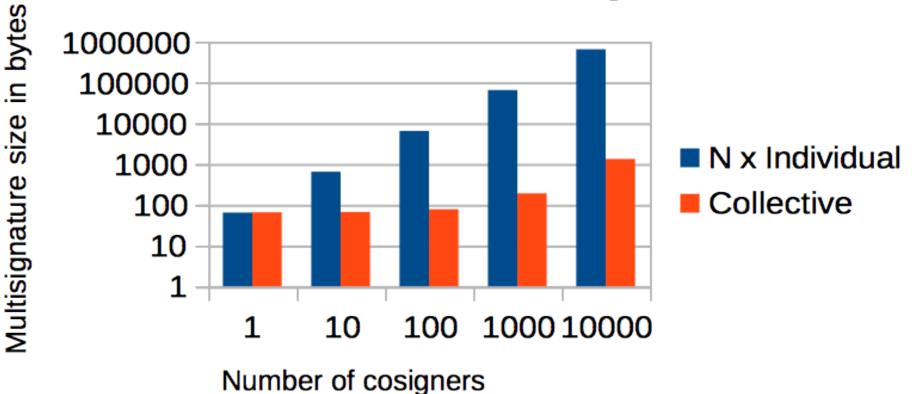
Collective versus individual signature verification



Results: Collective Signature Size

Ed25519: up to 512x smaller than N signatures

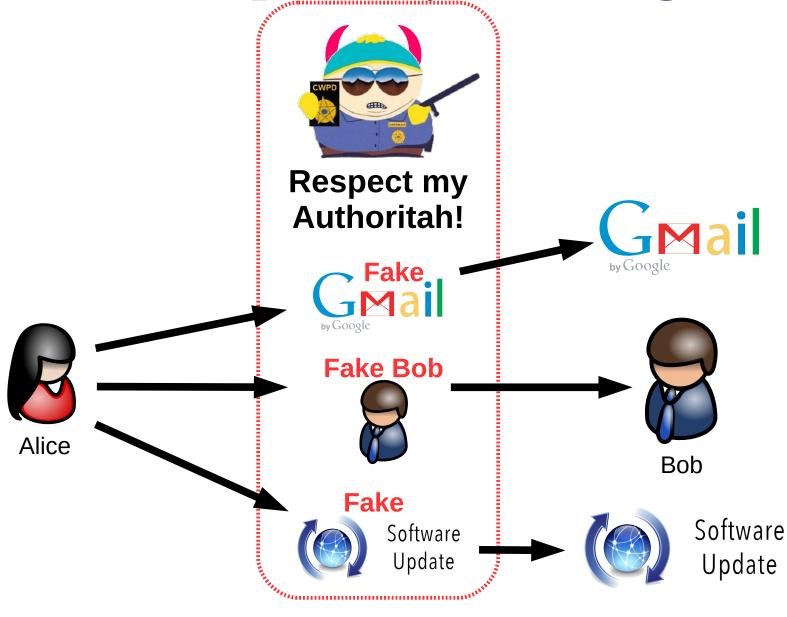
Collective versus individual signature size



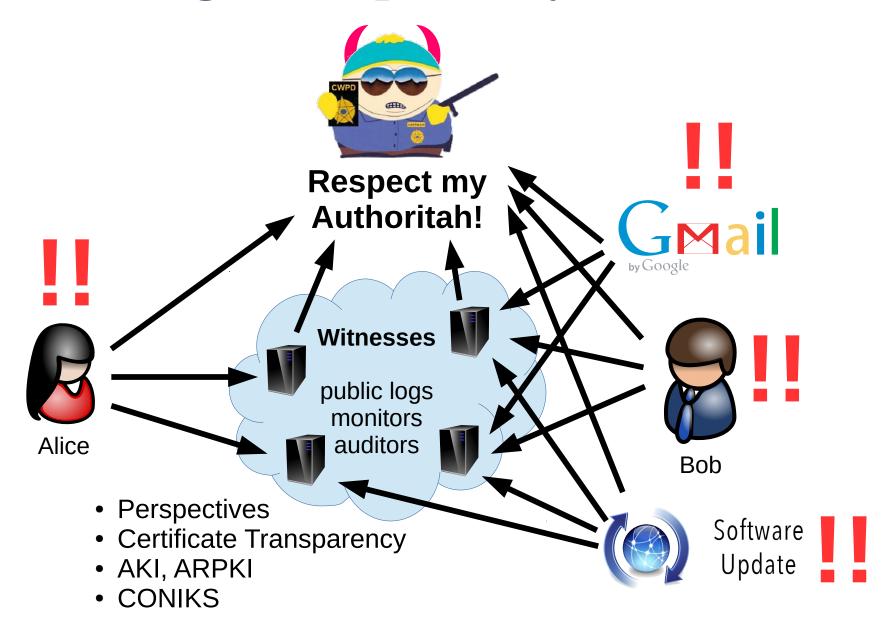
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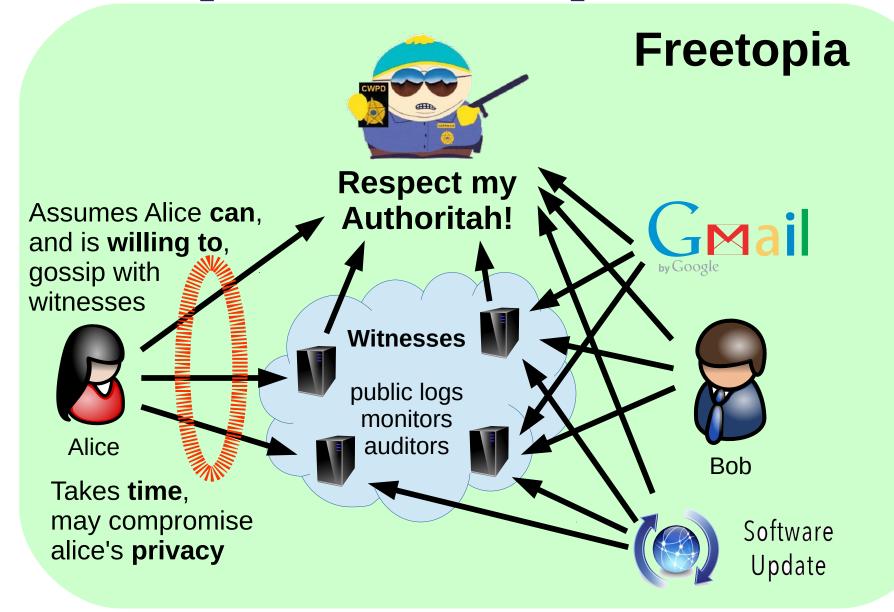
The Transparency Challenge



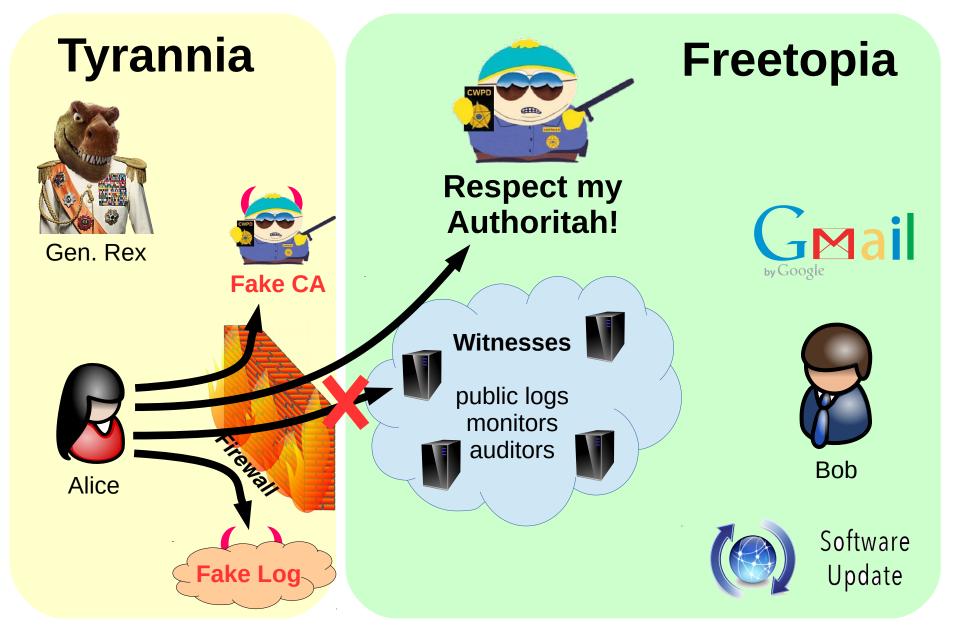
Existing Transparency Solutions



An Important Assumption



A Different Scenario



Gossip versus Collective Signing

Gossip can't protect Alice if she...

- Can't (because she's in Tyrannia)
- Doesn't want to (for privacy), or
- Doesn't have time to

cross-check each authoritative statements.

Collective signing **proactively** protects her from secret attacks even via her access network.

Attacker can't secretly produce valid signature

An "Extreme" Scenario

What if an attacker controls the target device, wants to secretly coerce the device's vendor to sign a back-doored operating system image?



- A phone sealed in a forensics lab can't gossip!
 - Certificate Transparency can't reveal its existence
- Only protection is to bind the transparency proactively into the device-verified signature

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Prototype available; give it a try!

Go to https://github.com/dedis/cosi

- Binaries: see releases
- Source: go get -u github.com/dedis/cosi

cosi sign -g group.toml -o sig msg_file cosi verify -g group.toml -s sig msg_file

Run your own witness server: cosi server Standalone verifiers for C, Go – see README

Status, Incremental Deployment

Still experimental! But...

- DEDIS lab committed to supporting, assisting with integration/deployment efforts
- Don't want to trust collective signatures yet?
 Add in extension field alongside individual sig
- Don't want to trust protocol, server liveness?
 Fork/exec 'cosi sign', set timer, kill if needed
- Don't want to trust cosi software?
 Sandbox it! Needs almost nothing to run.

Send feedback privately or discuss publicly on https://groups.google.com/forum/#!forum/cothority

Other uses of collective signing

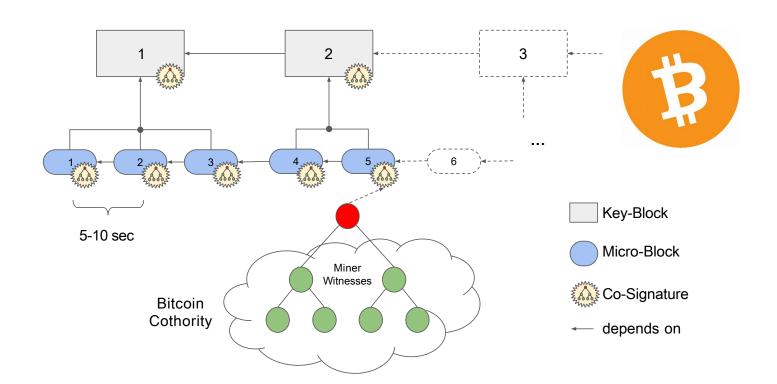


(credit: Tony Arcieri)

Other uses of collective signing

"Enhancing Bitcoin Security and Performance with Strong Consistency via Collective Signing"

- To appear at USENIX Security 2016
- Draft: http://arxiv.org/abs/1602.06997



Conclusion

Grand challenge: decentralize all the authorities!

Practical baby step: decentralized witness cosigning

- Ensures that for any signed statements that exists,
 many parties have witnessed, publicly logged it
 - Protects even relying parties that can't gossip
- Can incrementally add to existing authorities
- CoSi protocol scales to large witness groups

Available: https://github.com/dedis/cosi

Public question/answer, discussion forum: https://groups.google.com/forum/#!forum/cothority

Scalable Collective Timestamping

Like classic **digital timestamp** services, only decentralized.



- Each round (e.g., 10 secs):
 - 1) Each server collects hashes, nonces to timestamp
 - 2) Each server aggregates hashes into Merkle tree
 - 3) Servers aggregate local trees into one global tree
 - 4) Servers collectively sign root of global tree
 - 5) Server give signed root + inclusion proof to clients
- Clients verify signature + Merkle inclusion proof

Verifiably Fresh Software Updates

Alice accepts only updates with fresh timestamp:

- Knows update can't be an outdated version: tree contains inclusion proof of her nonce
- Knows update can't have targeted backdoor: witness cothority ensures many parties saw it

