# Resource-Aware Session Types for Digital Contracts

#### **Ankush Das**

Stephanie Balzer

Jan Hoffmann

Frank Pfenning

Ishani Santurkar

Carnegie Mellon University

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Digital Contracts



**Digital Contracts** 



### Challenges:

- 1. contract protocols are violated
- 2. asset duplication/deletion
- 3. automatic inference of execution cost



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Programs implementing a digital (legal or financial) transaction

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**Online Transactions** 

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Wealth Management Applications

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**Online Transactions** 

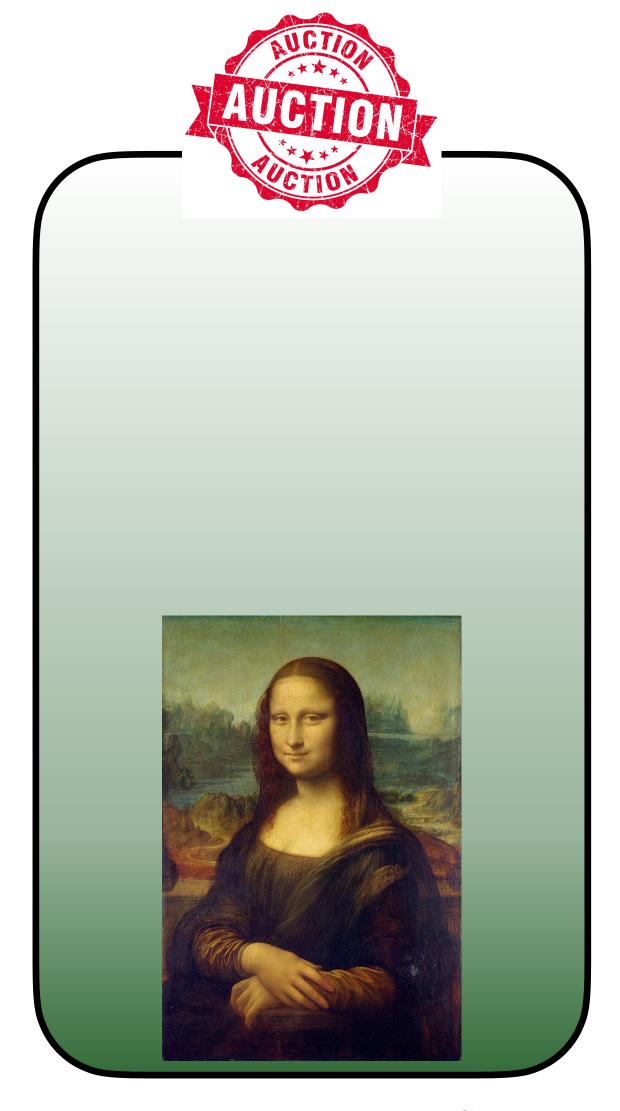


Wealth Management Applications

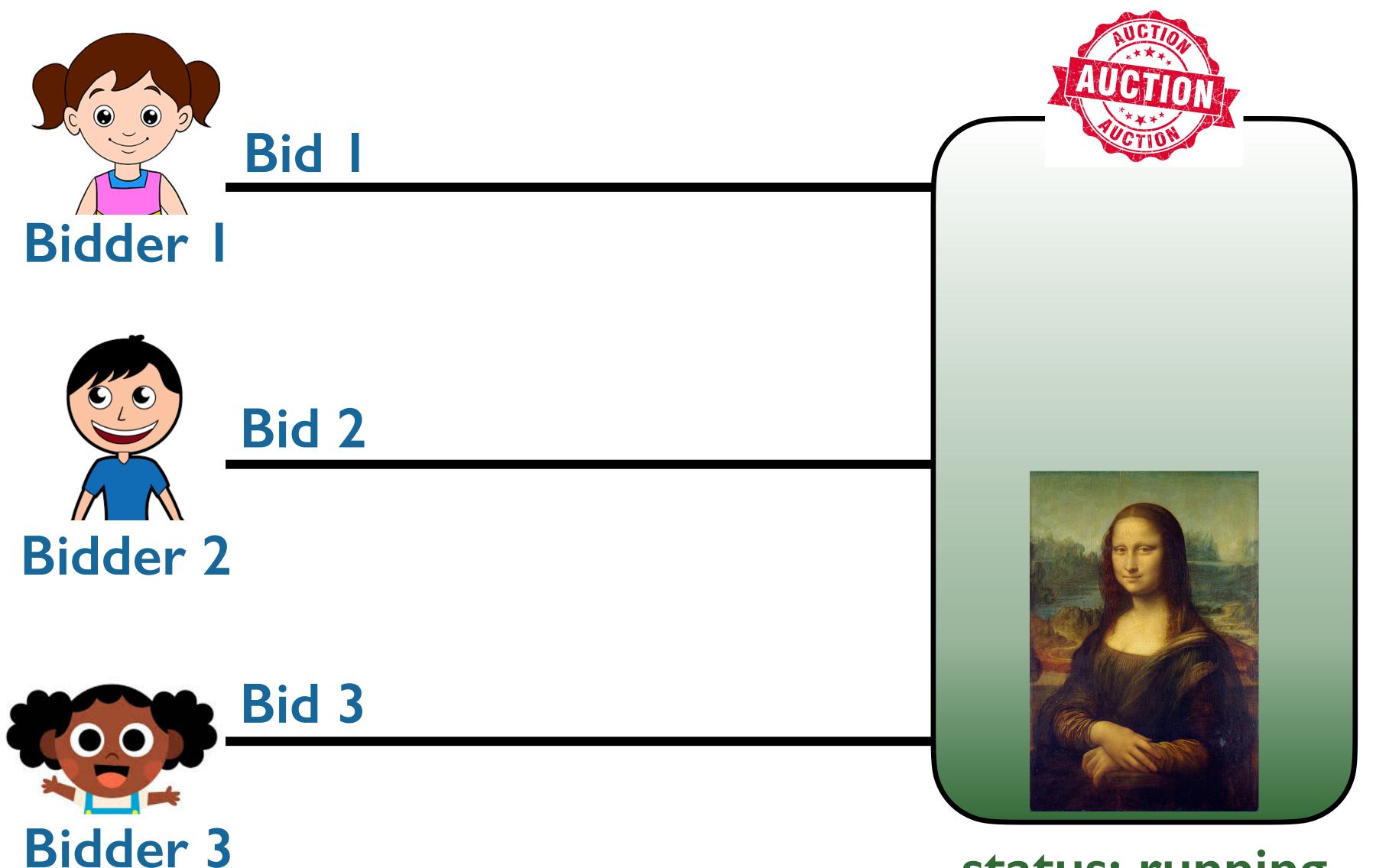
recently popularized in the form of



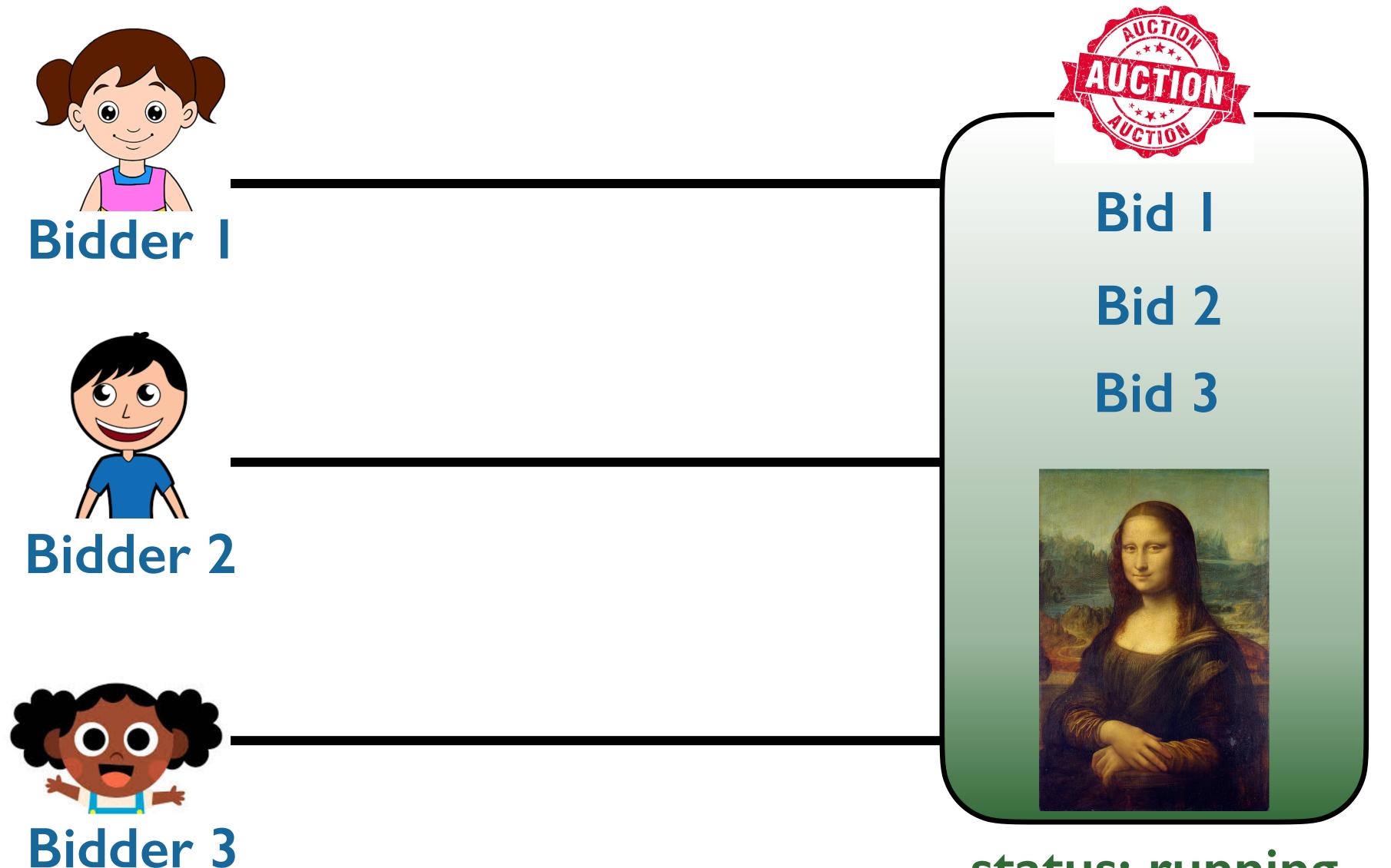
Blockchains and Cryptocurrencies

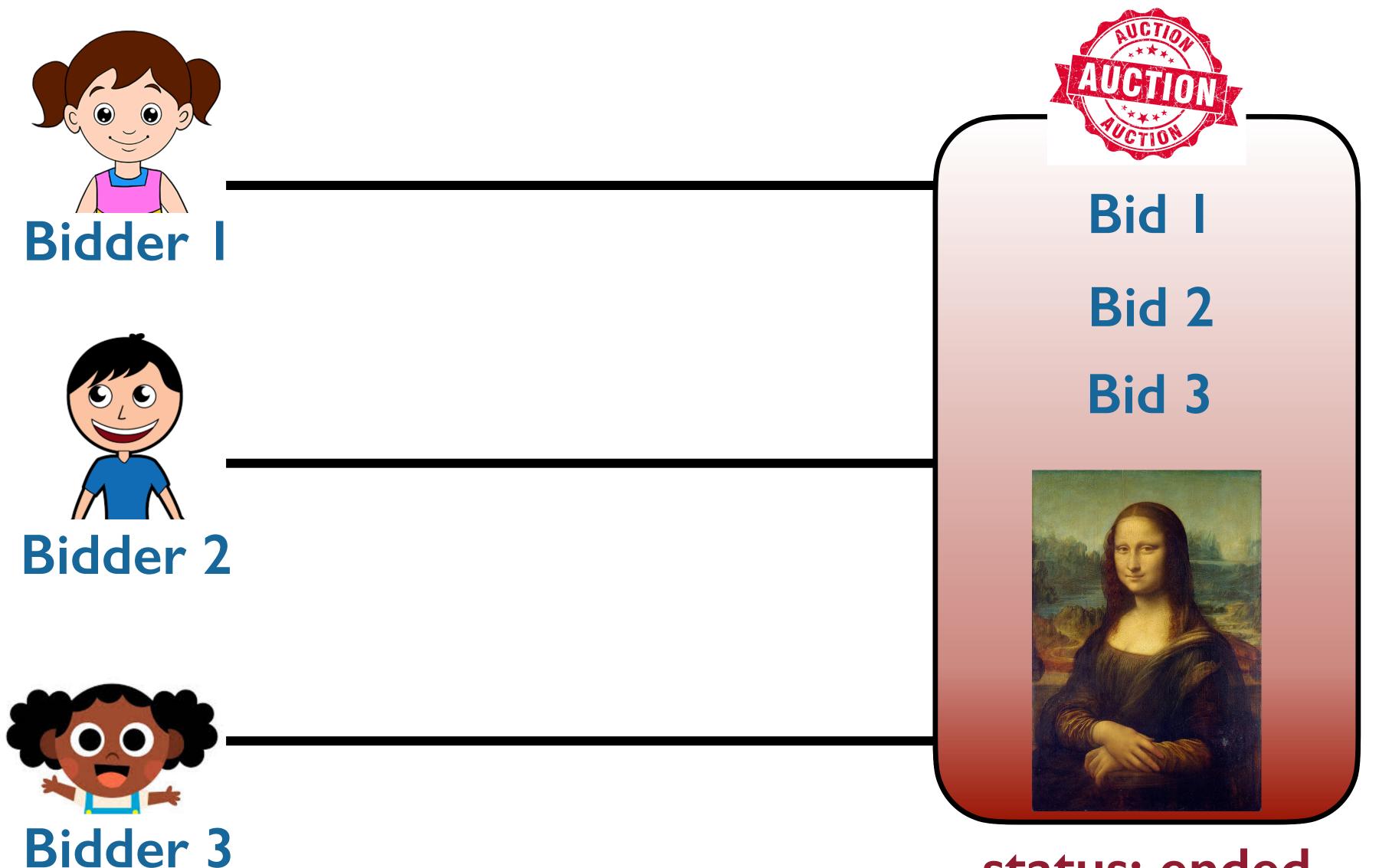


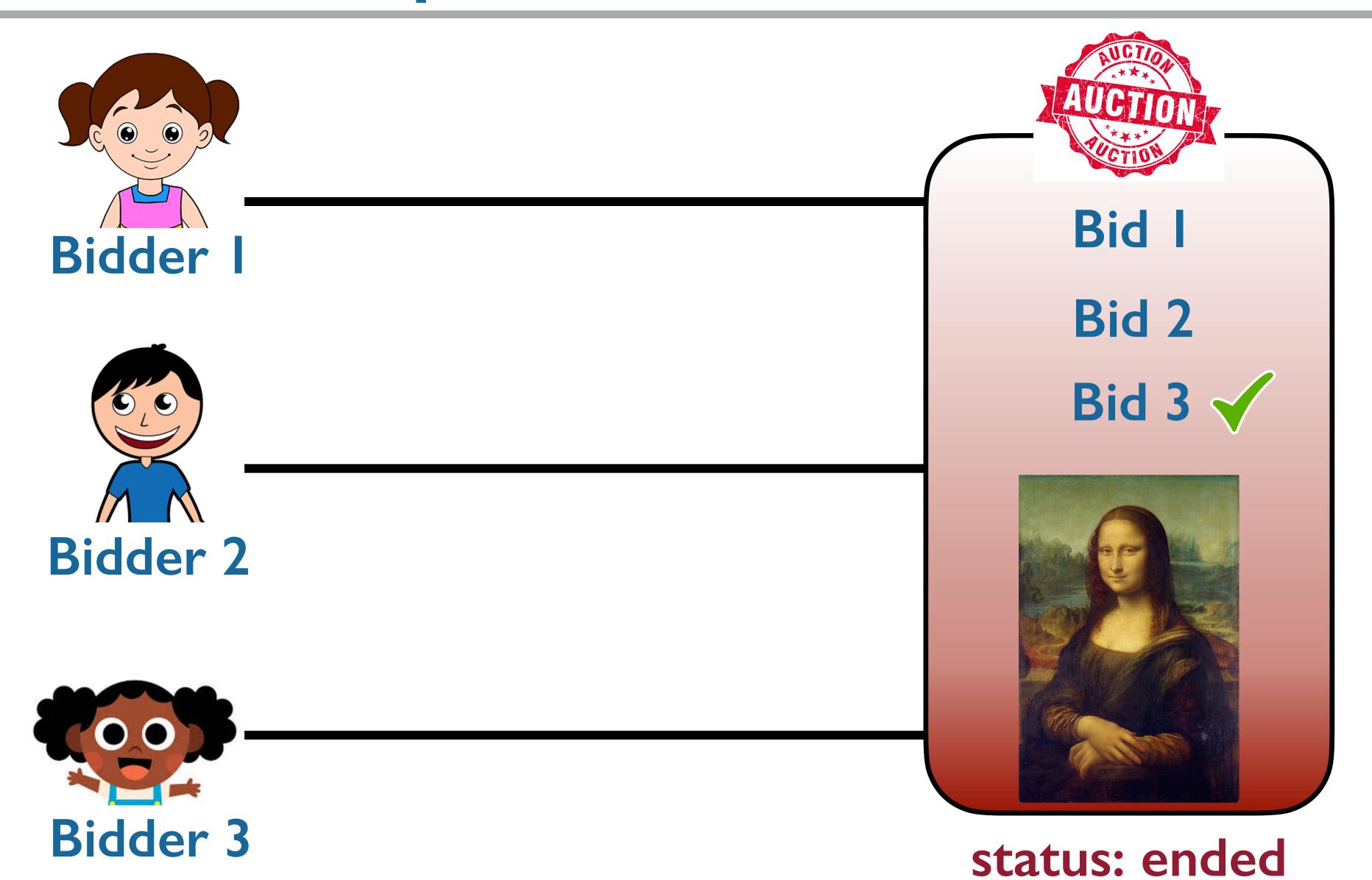
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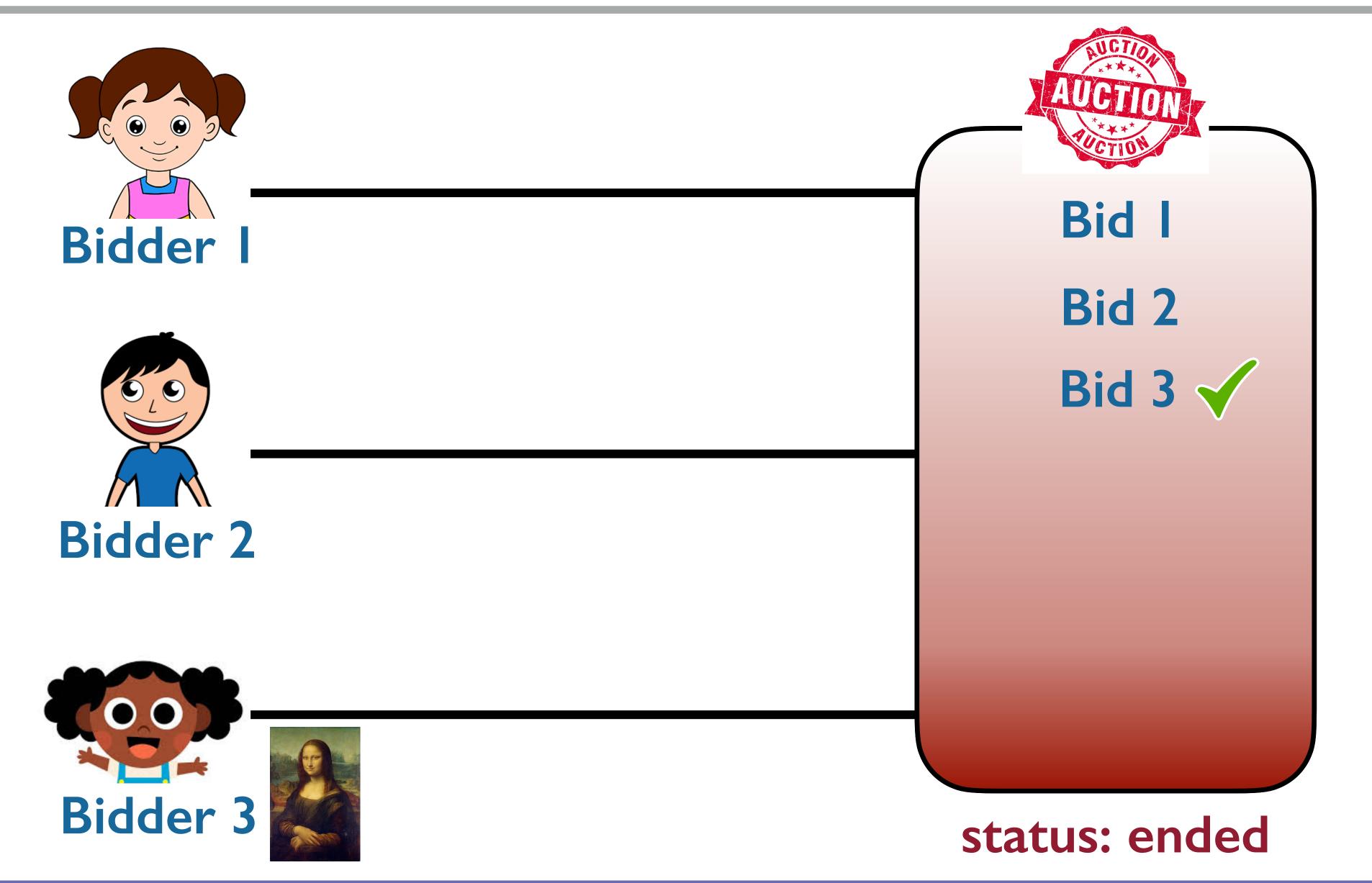


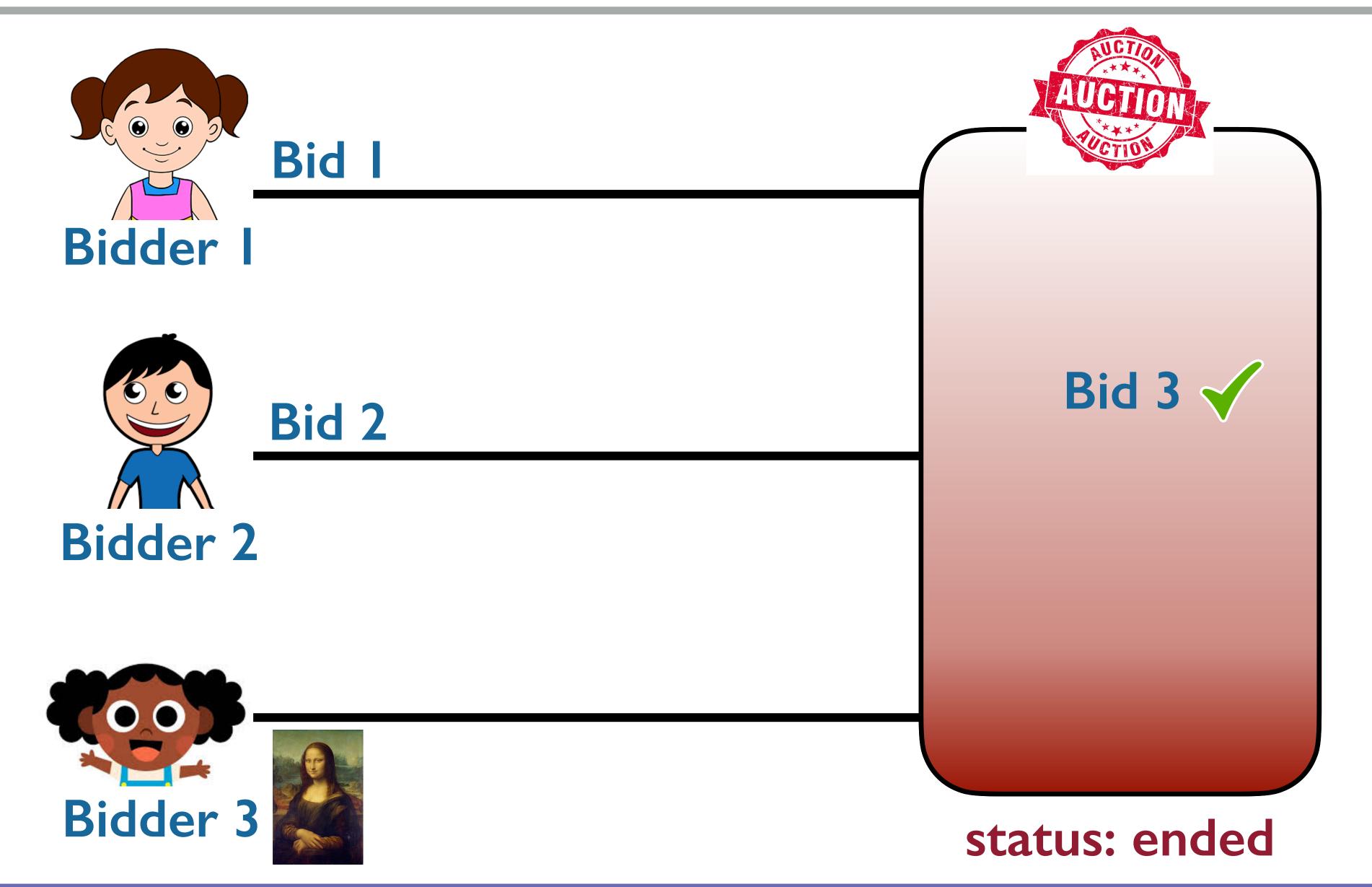
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# Auction in Solidity \$\frac{1}{2}

```
function bid() public payable {
  bid = msg.value ;
  bidder = msg.sender ;
  pendingReturns[bidder] = bid ;
  if (bid > highestBid) {
    highestBidder = bidder ;
    highestBid = bid ;
function collect() public returns (bool) {
  require (msg.sender != highestBidder) ;
  uint amount = pendingReturns[msg.sender] ;
 msg.sender.send(amount);
  return true ;
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What happens if collect is called when auction is running?

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# Auction in Solidity 5

What happens if collect is called when auction is running?

add clause:

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How much fee should the user pay? No support for static and automatic cost prediction!



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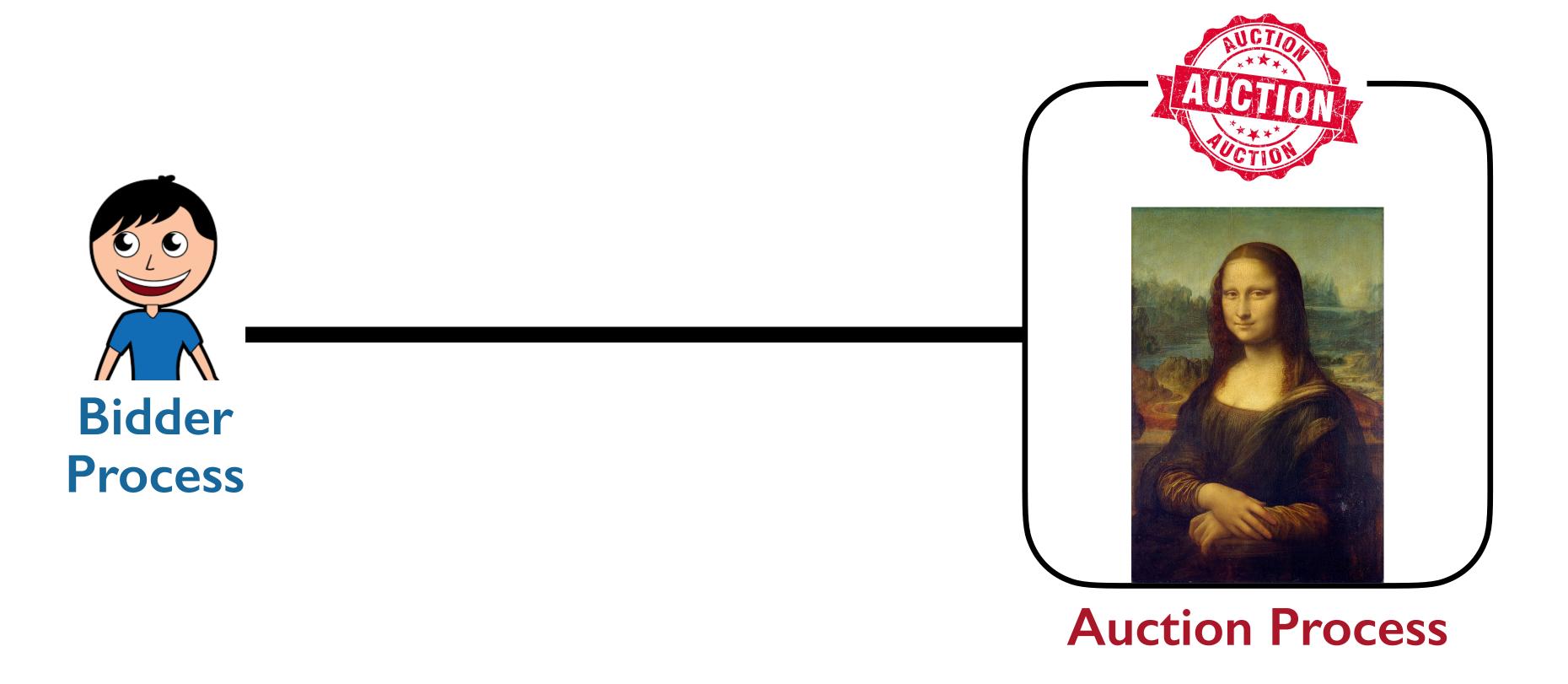


#### **Solutions:**

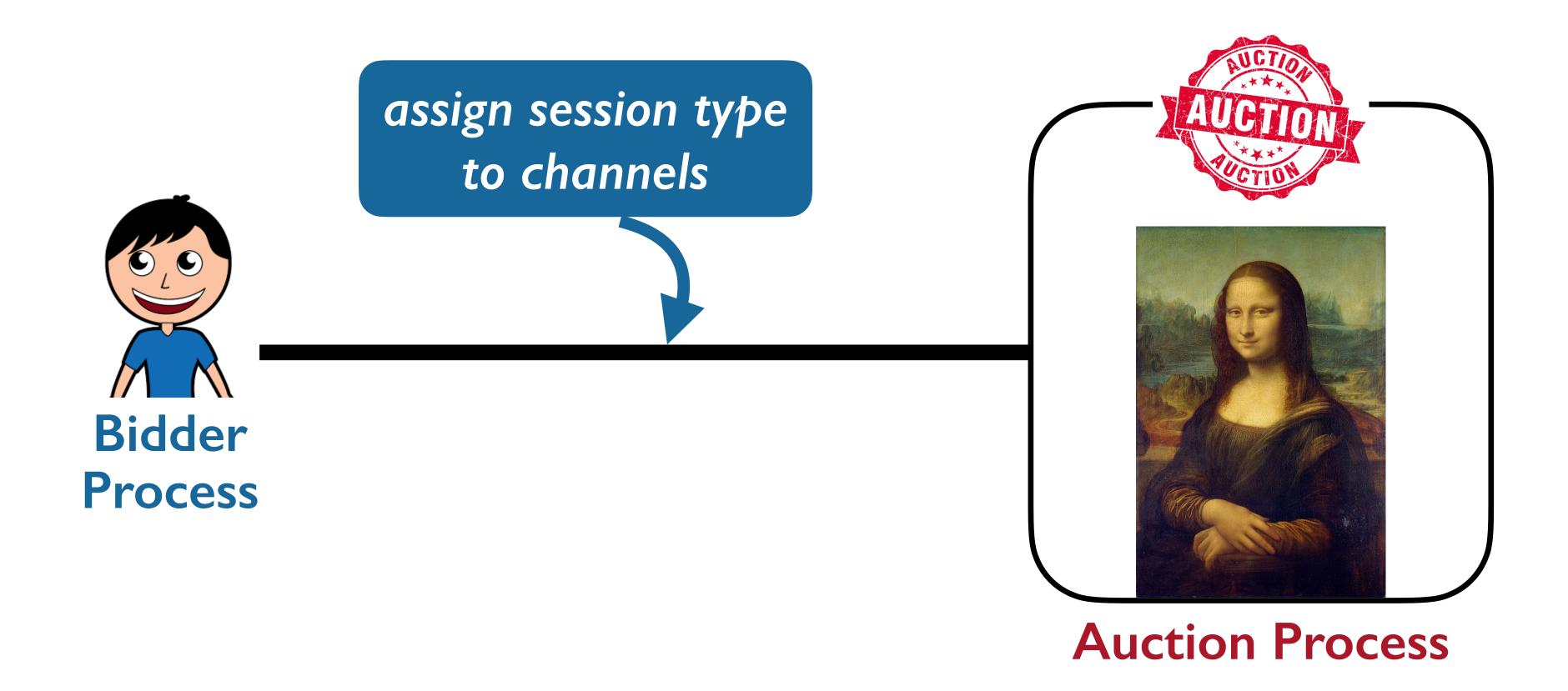
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Type system for implementing concurrent message-passing processes

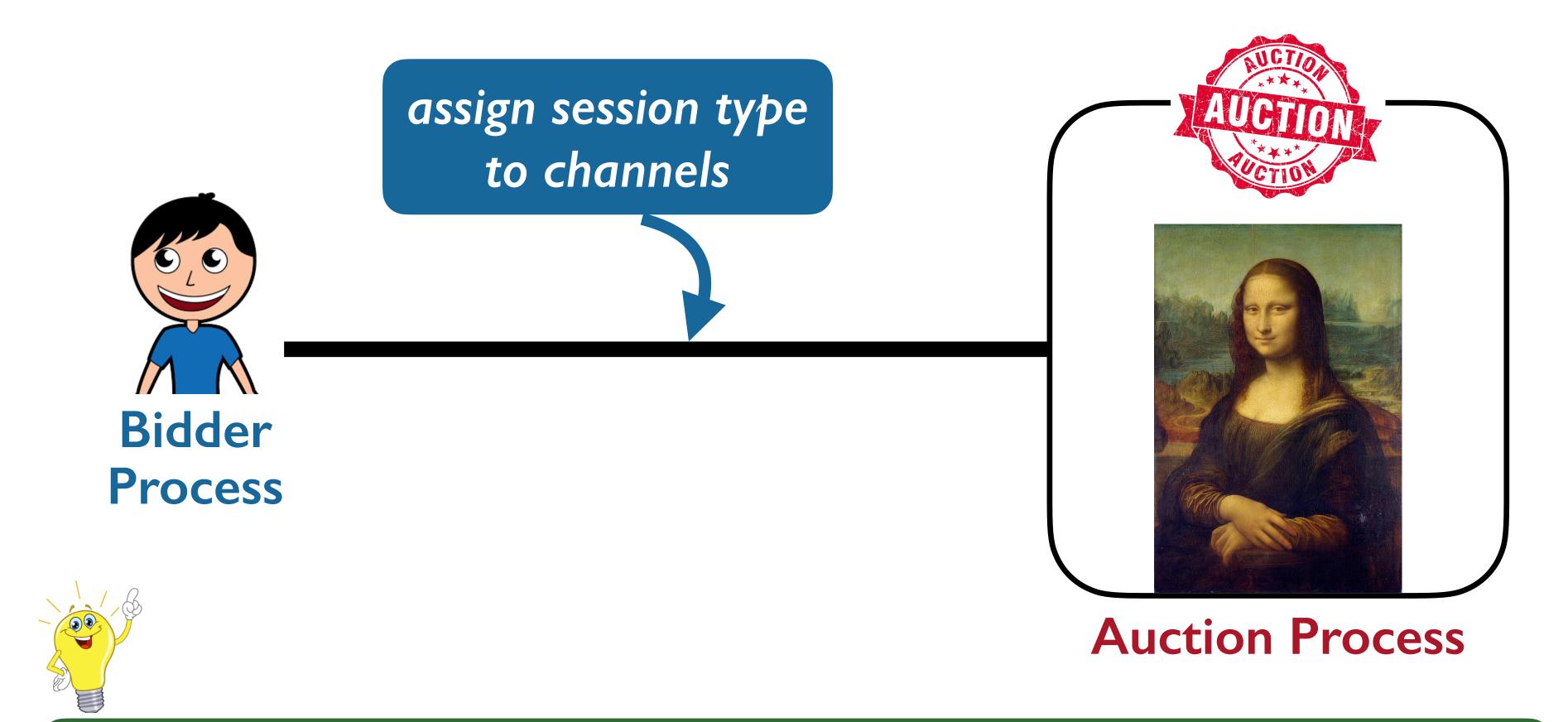
Type system for implementing concurrent message-passing processes



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Key Idea: Nomos uses session types to express contract protocols

```
type \ \mathsf{auction} = \oplus \{\mathbf{running} : \& \{\mathbf{bid} : \mathsf{money} \multimap \mathsf{auction}\}, \\ \mathbf{ended} : \& \{\mathbf{collect} : \oplus \{\mathbf{won} : \mathsf{monalisa} \otimes \mathsf{auction}, \\ \mathbf{lost} : \mathsf{money} \otimes \mathsf{auction}\}\}\}
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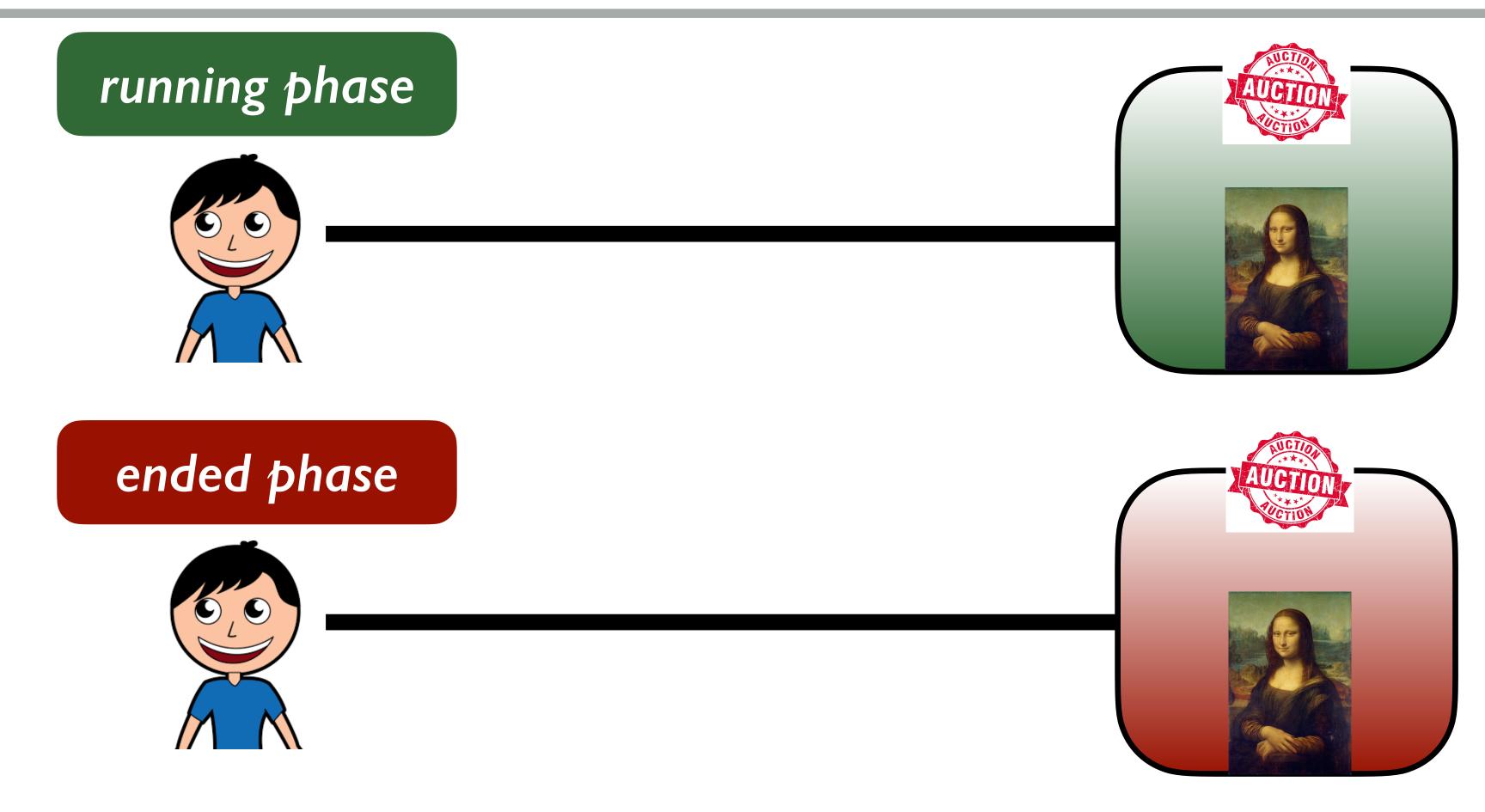
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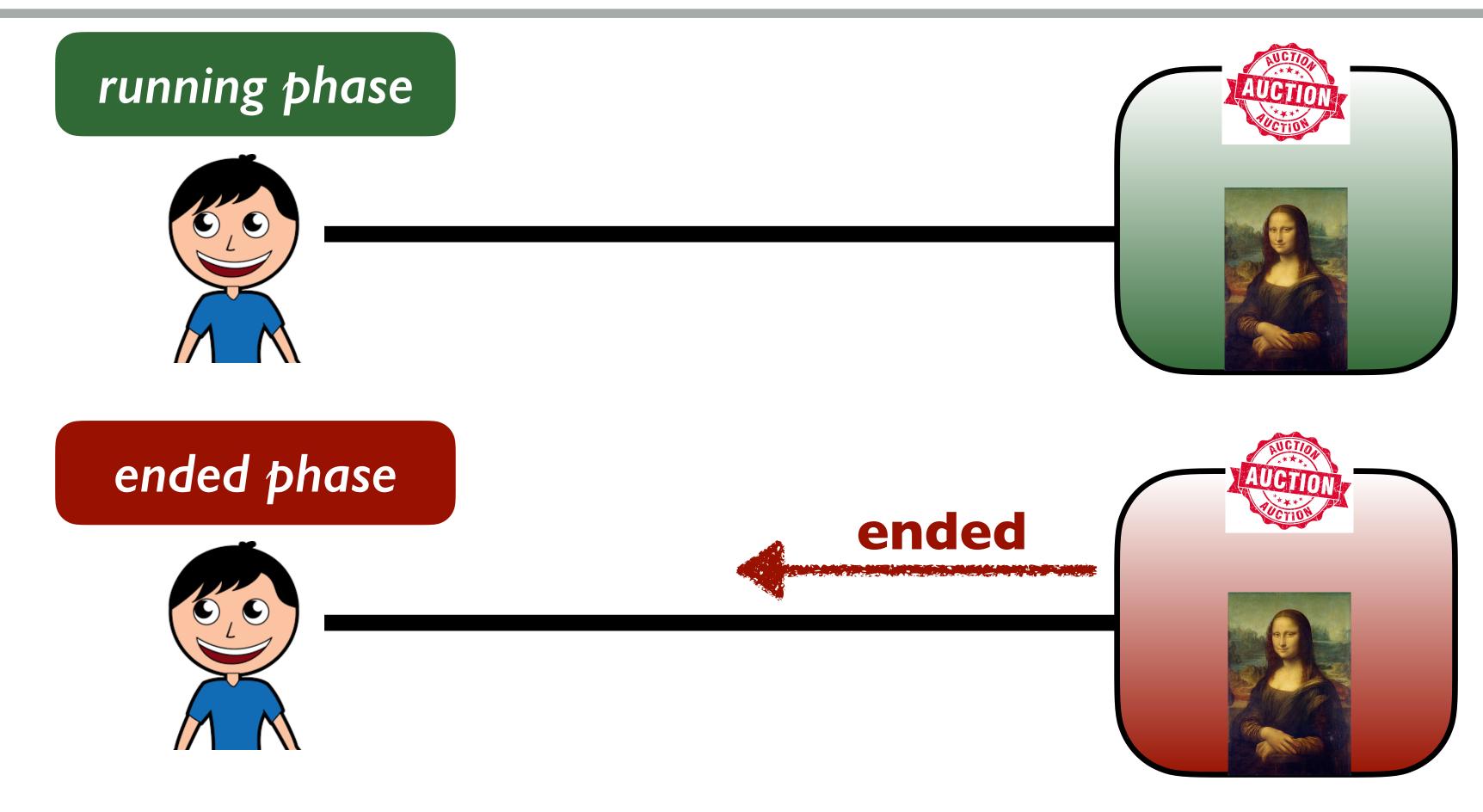
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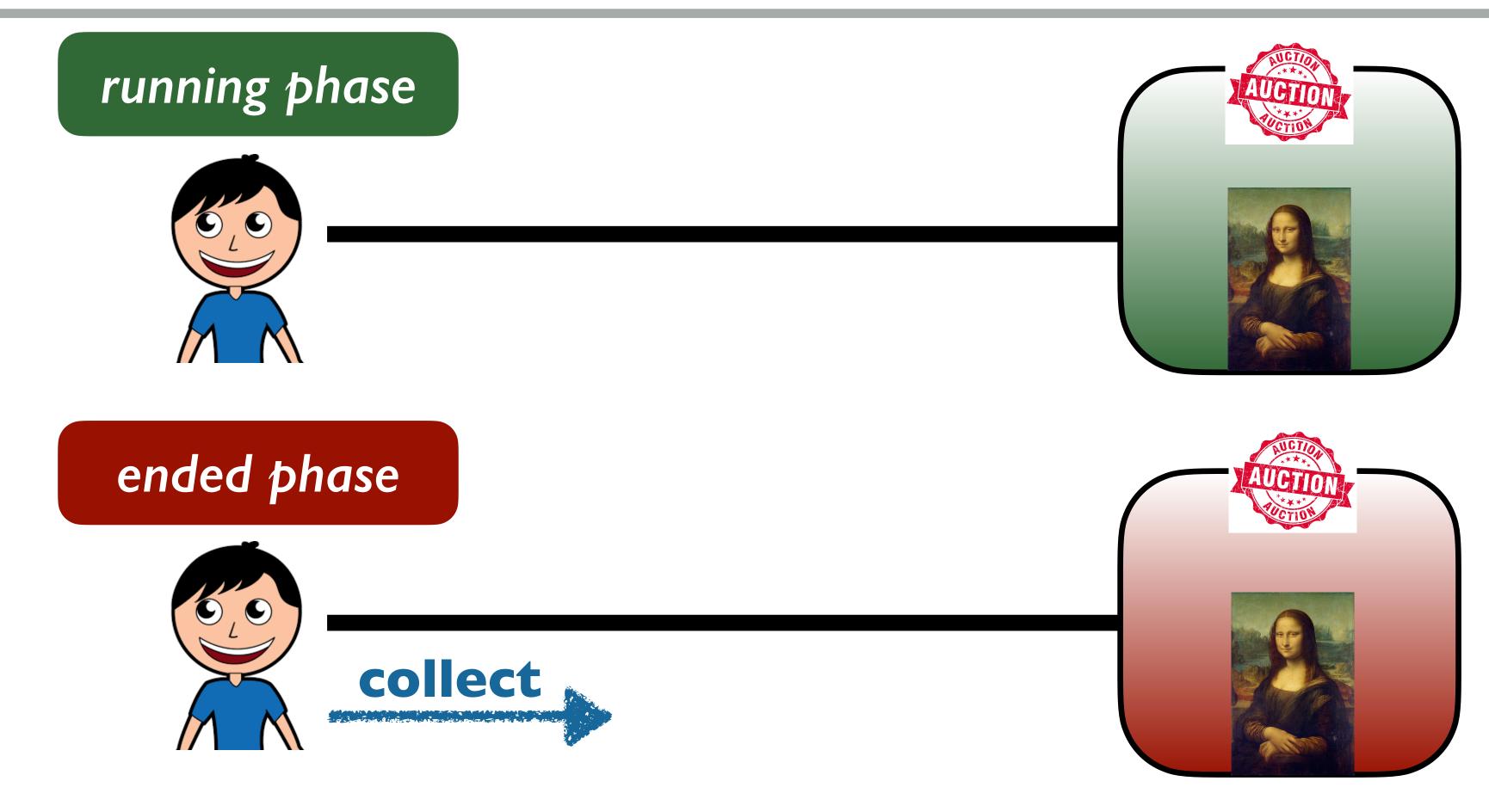
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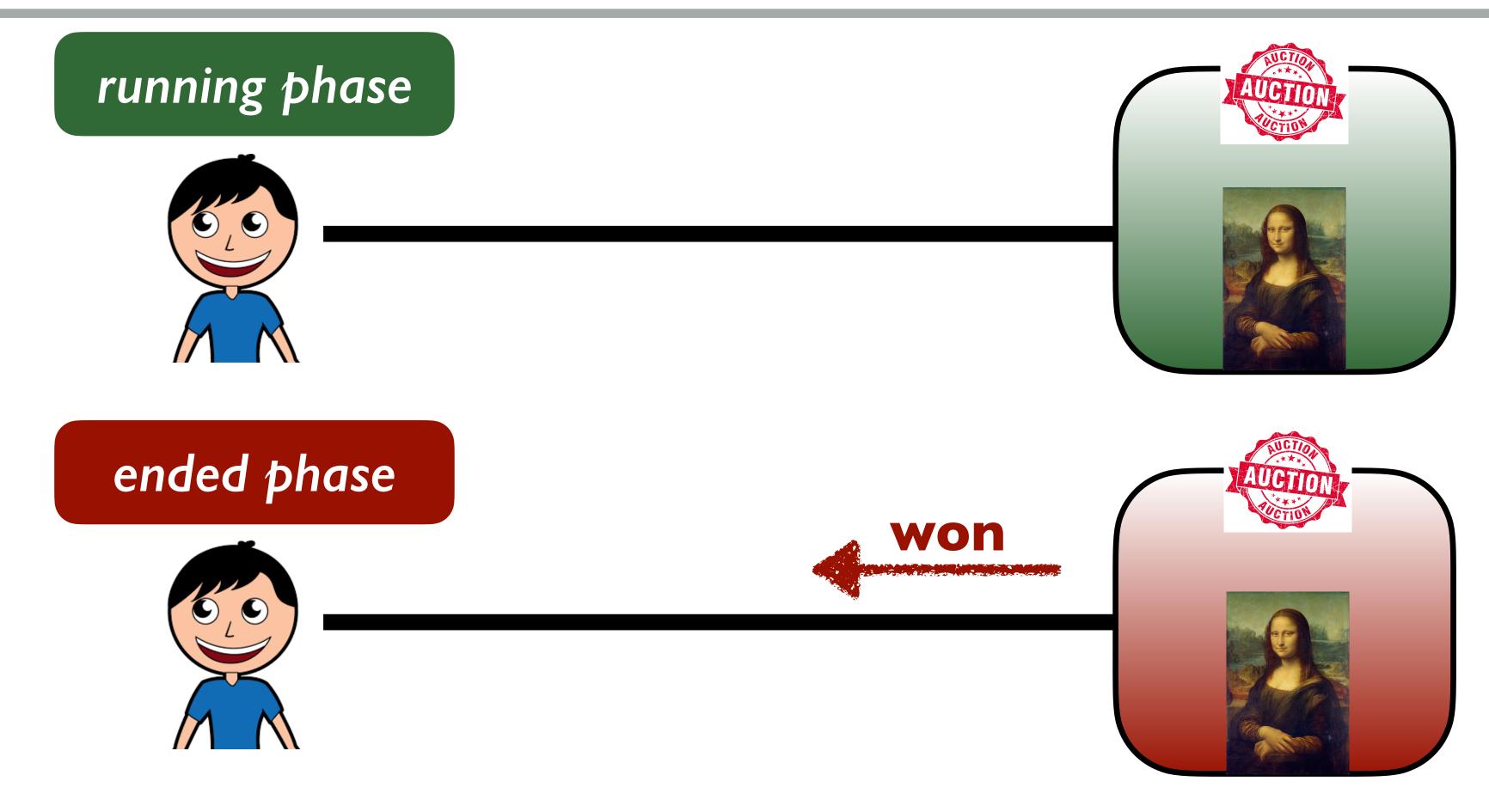
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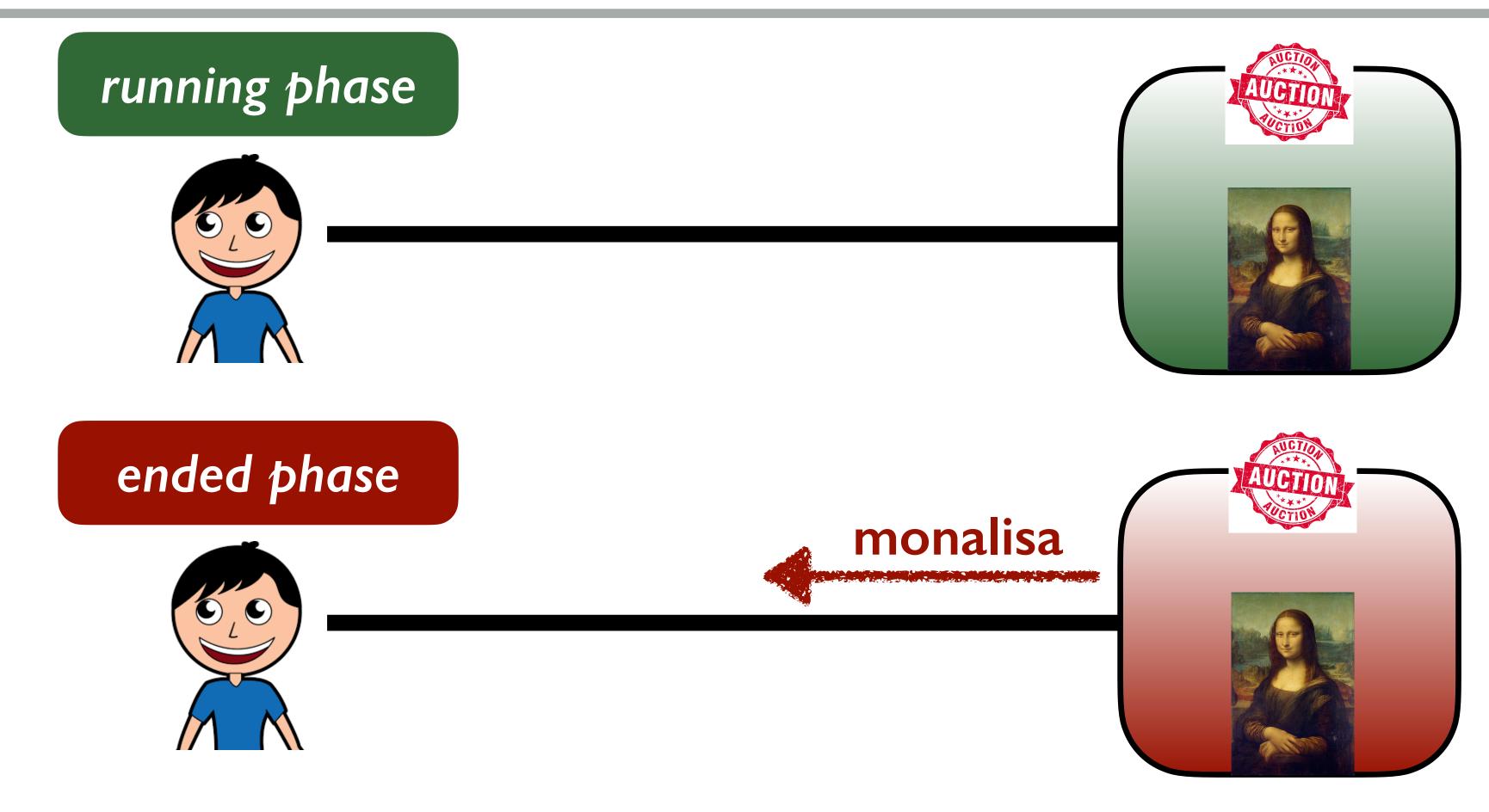
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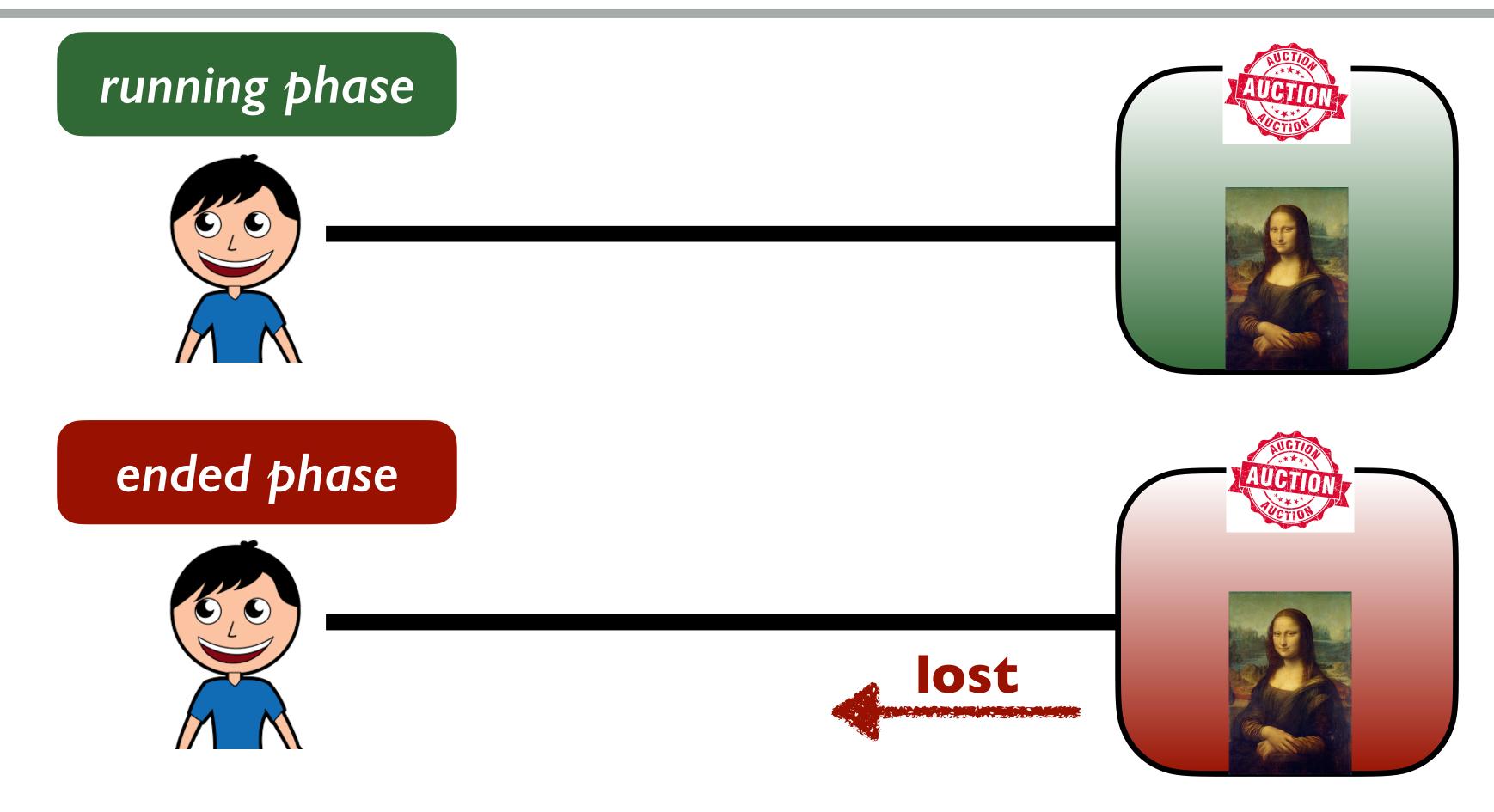
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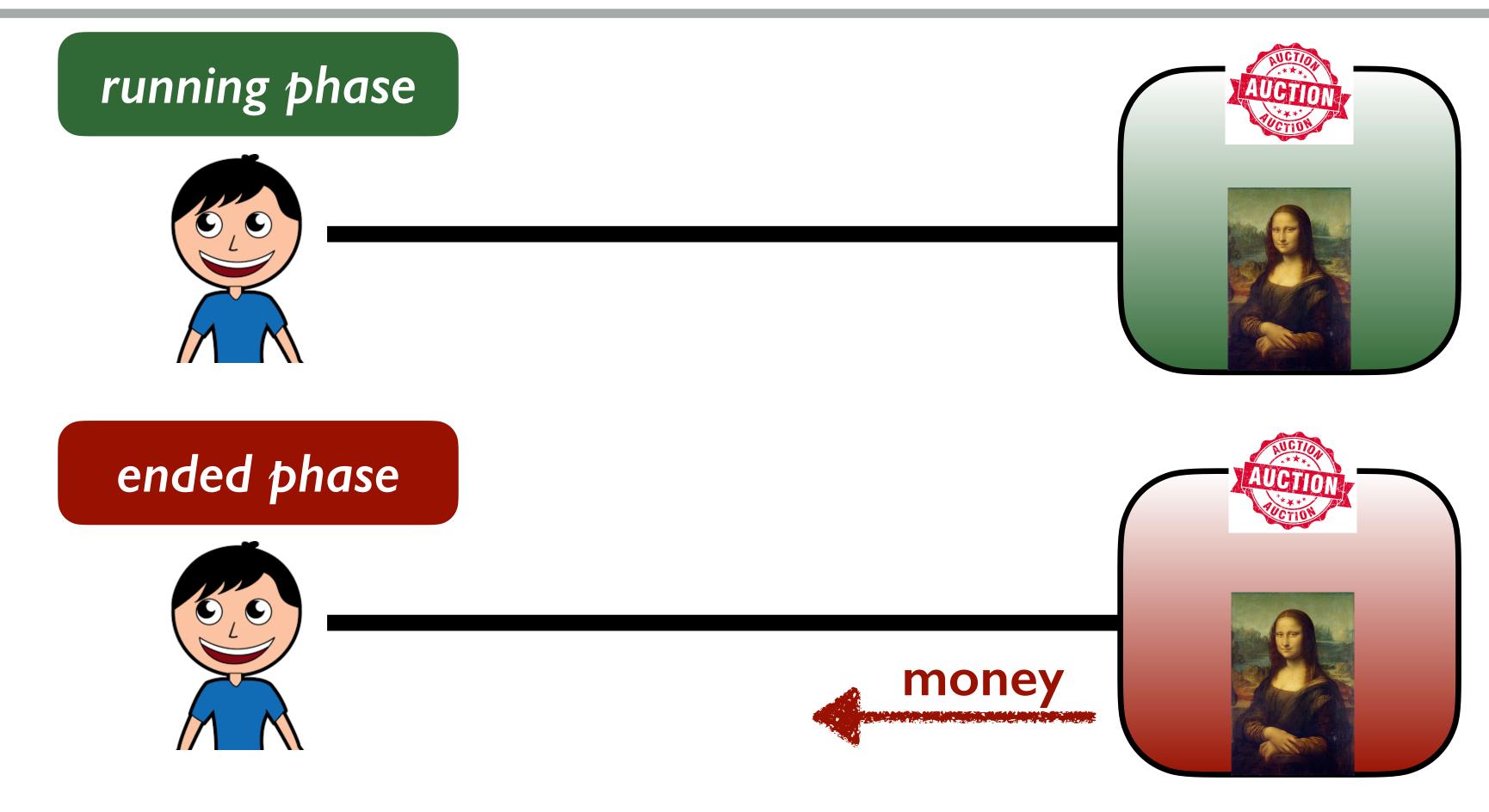
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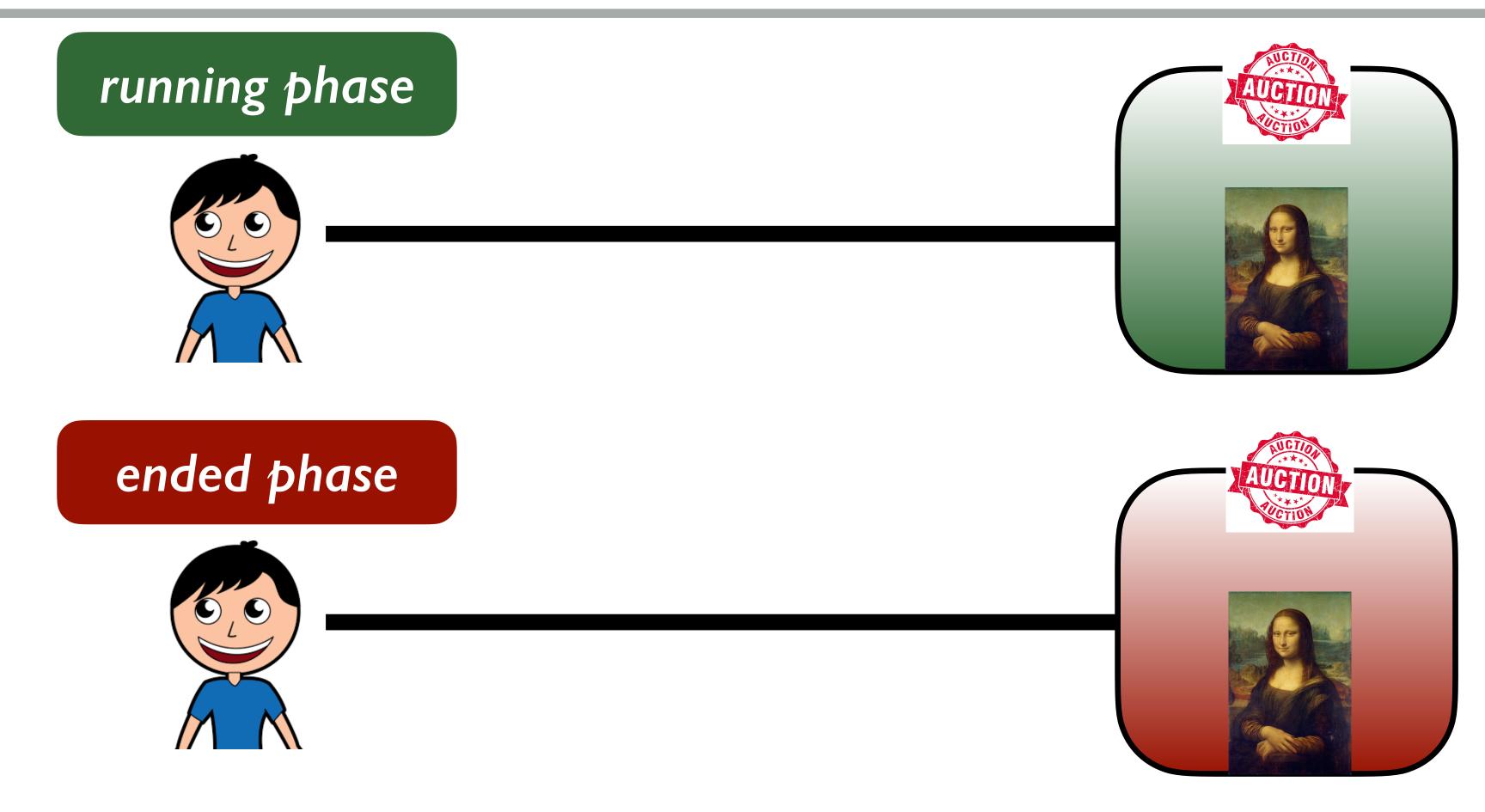
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#### Nomos statically guarantees:

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### Get Asset Preservation for Free!



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user pays 3 units as fees at start of execution

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Potential annotations are automatically inferred using LP solver

## Conclusion: Features of Nomos

- > session types enforce contract protocols
- Inear types track assets, prevent duplication/deletion
- resource-aware types infer execution cost
- re-entrancy attacks eliminated by construction
- type safety theorem guarantees all properties
- extensive evaluation on standard digital contracts