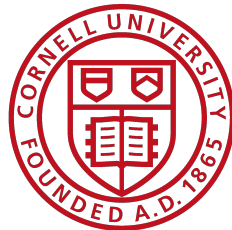


Towards Automated Safety Vetting of PLC Code in Real-World Plants

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Safety Hazards are Unique Threats in ICS

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News about Cyberattacks on Iran — Stuxnet and Flame, including commentary and archival articles published in The New York Times.

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July 6, 2017

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CBS NEWS / June 23, 2017, 12:48 PM

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The Russian attacks on Ukraine's power grid were extensive. In 2015, electricity was cut to nearly a quarter-million Ukrainians, and about a year later a transmission station was taken down, revealing the attacks were becoming more sophisticated.

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Hack attack causes 'massive damage' at steel works

22 December 2014 | Technology



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The hack attack led to failures in plant equipment and forced the fast shut down of a furnace

A blast furnace at a German steel mill suffered "massive damage" following a cyber attack on the plant's network, says a report.

PLC being a Major Attack Vector

```
21 IF Pallet_Sensor AND NOT(Part_Sensor) THEN
22   Pallet_Arrival := true;
23 END_IF;
24
25 IF Part_Sensor THEN
26   Retire_Pallet := true;
27 END_IF;
28
29 IF Pallet_Arrival AND CNC_Part_Process AND Robot_Ready AND
30   NOT(Part_Sensor) THEN
31   Update_Part_Process := true;
32   CNC_Part_Ready := false;
33   Robot_Ready := false;
34 END_IF;
```

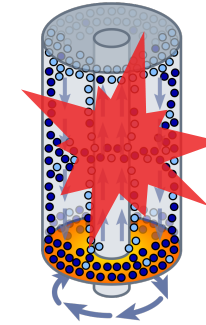
**Controller Code w/
Safety Violations**

Insider Attacks or Bugs



**Programmable
Logic Controller
(PLC)**

**Core Control Unit on
the Factory Floor**

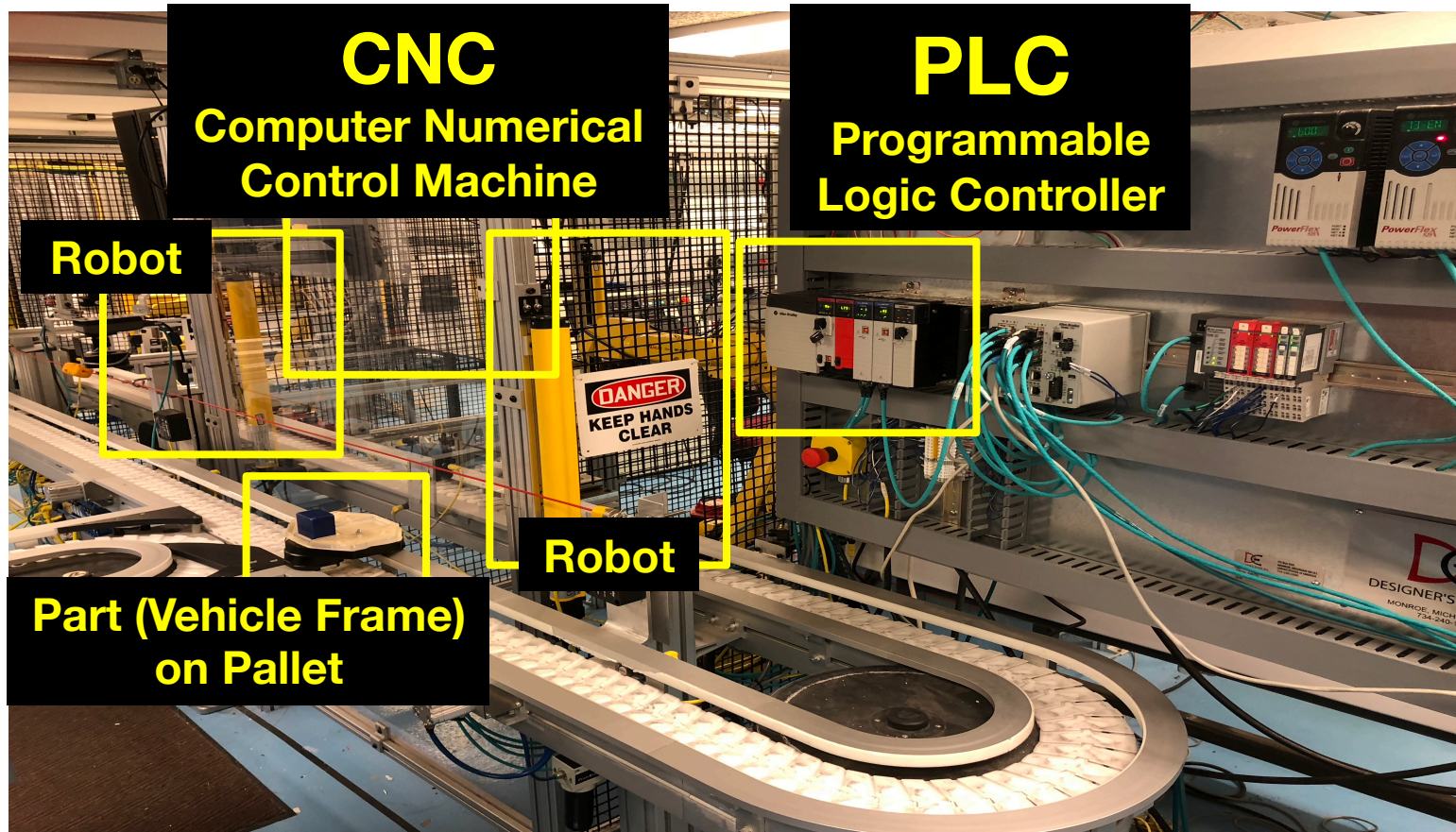


Physical Damage

**Different from Financial
Loss Often Seen in Attacks
in Consumer Systems**

A great many of prior work: e.g., TSV (NDSS'14), SymPLC (FSE'17)

Overlooked Fact: ICS is Complex; PLC is NOT Working Alone



Real-world Automotive Manufacturing Testbed

Developed by **No.1 Vendor**
(Rockwell Automation)

PLCs are **driven by events**
from other machines

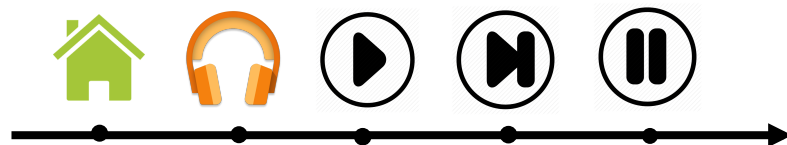
Testing PLC code requires external event inputs

Testing Event-driven Code in Other Domains

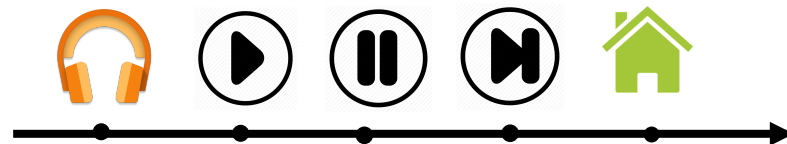
– Simulating and Rearranging Events

Android App: Anand *FSE'12*, Jensen *ISSTA'13*,
Mirzaei *Softw. Eng. Notes'12*, Yang *CCS'13*

Web Program: SymJS *FSE'14*, Saxena *Oakland'10*

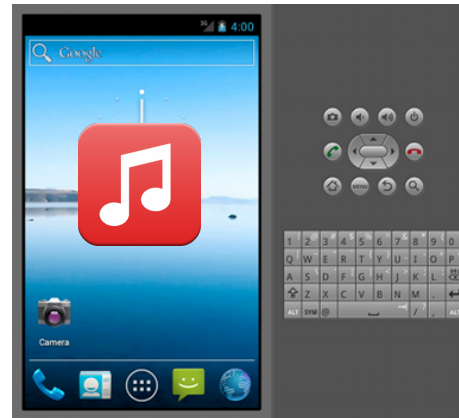


Simulated Event Sequence



Rearrange Event Order

⋮

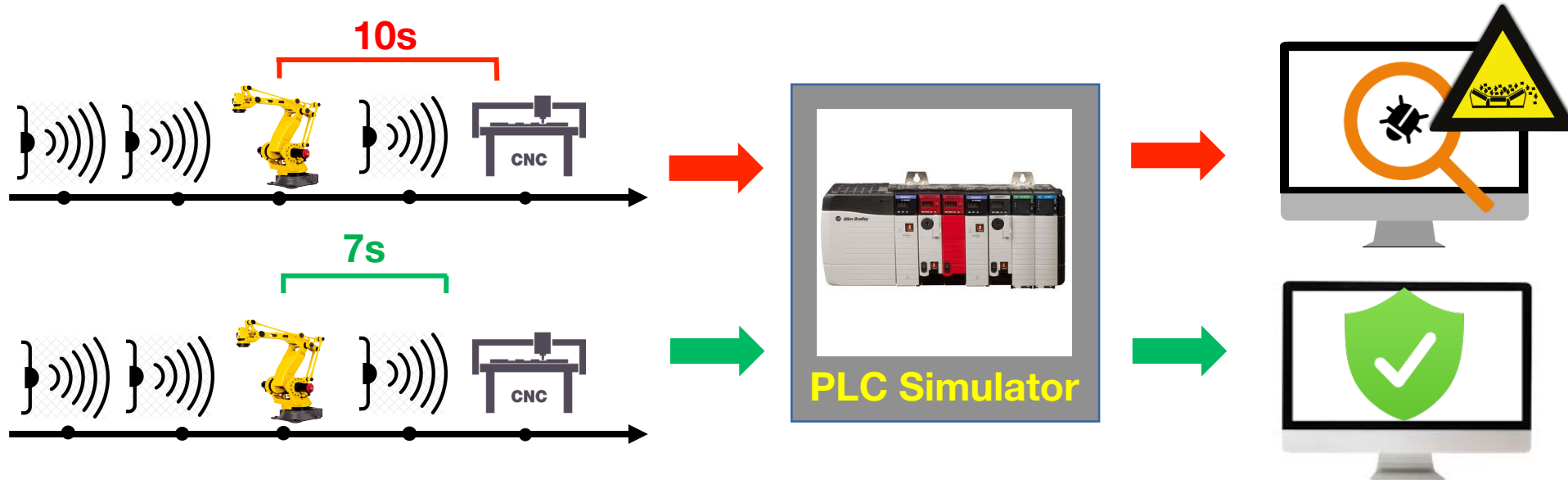


**App Testing in
Emulator**



Crash

Rearranging Event Order to Test PLC Code is **NOT Sufficient**



Event Sequences of Same Ordering **But Different Timings**

Timing factor: Nature of ICS

Timeliness, Throughput



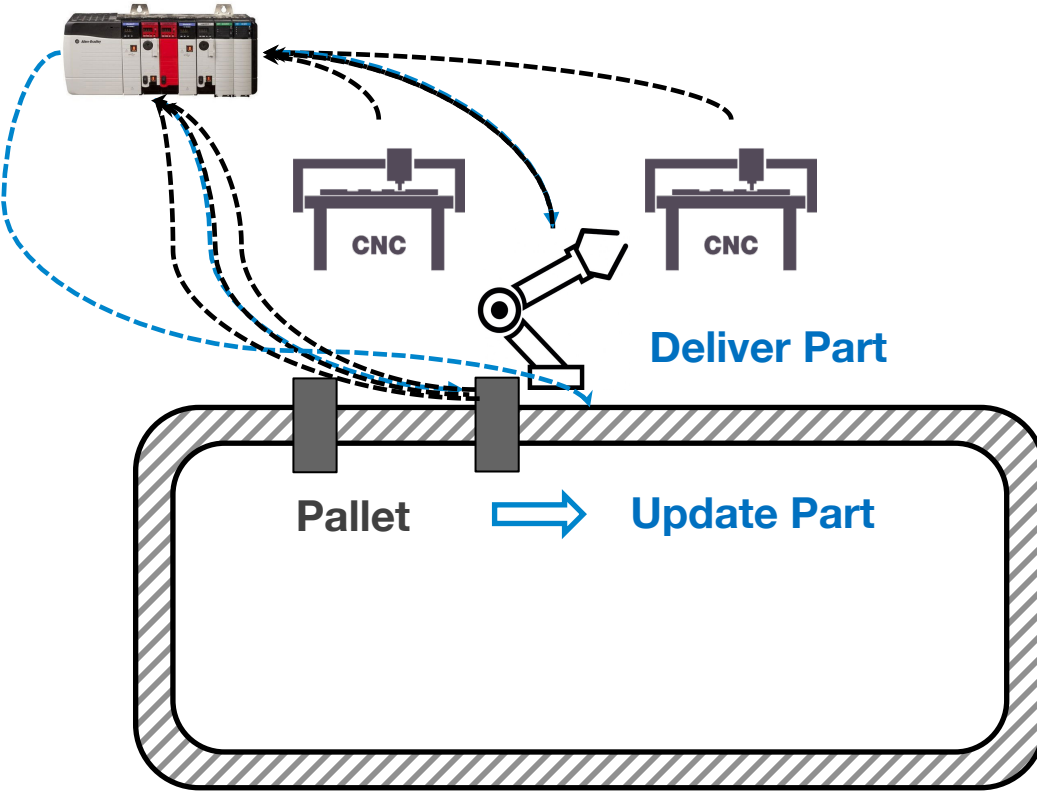
Internal Timeouts

Machine Speed Limits

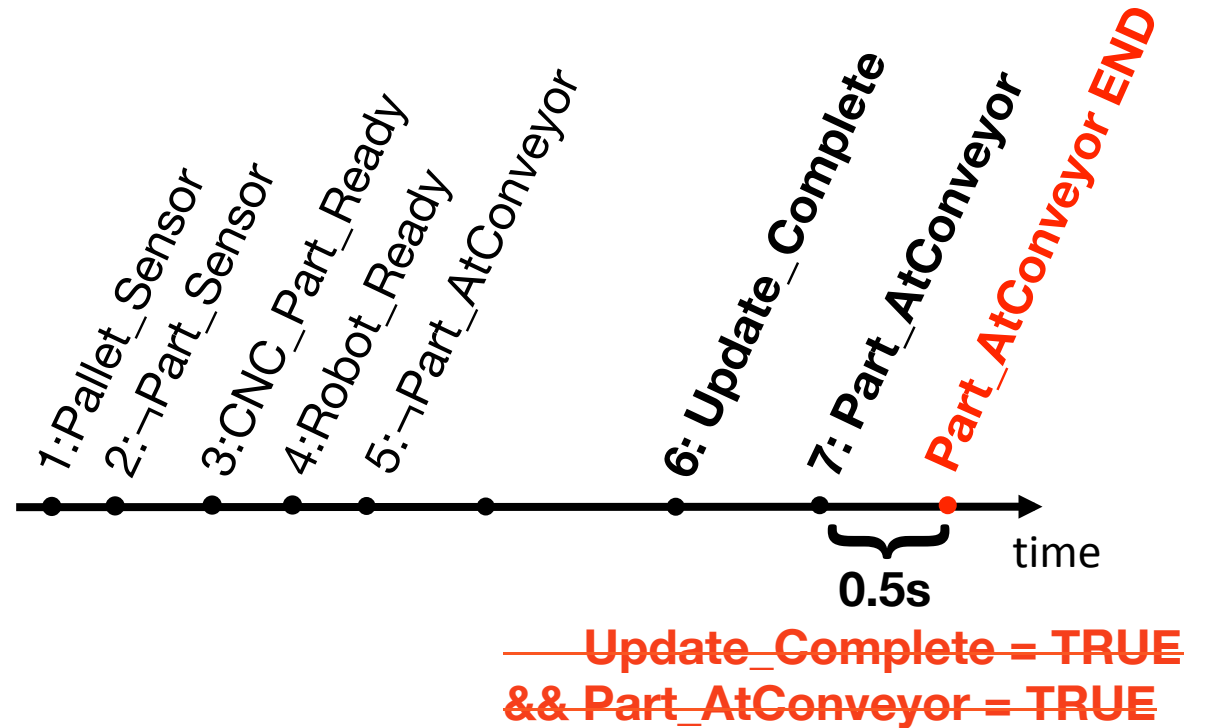


External Timing Constraints

A Running Example



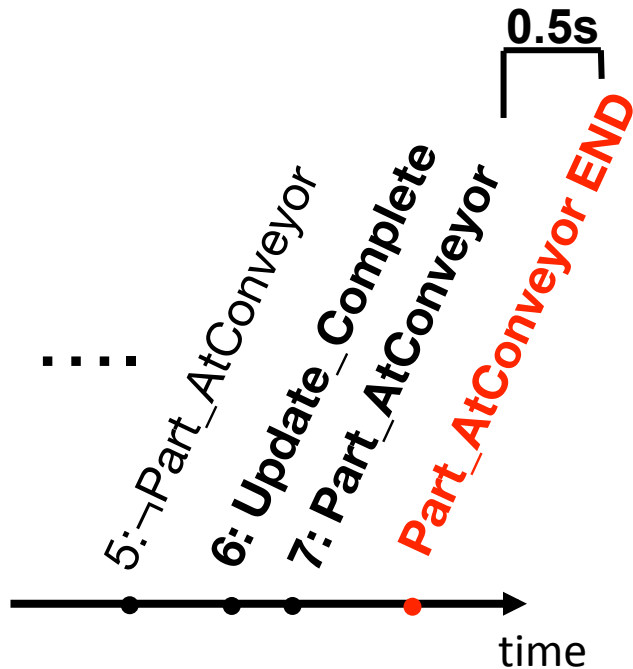
Events Received by PLC



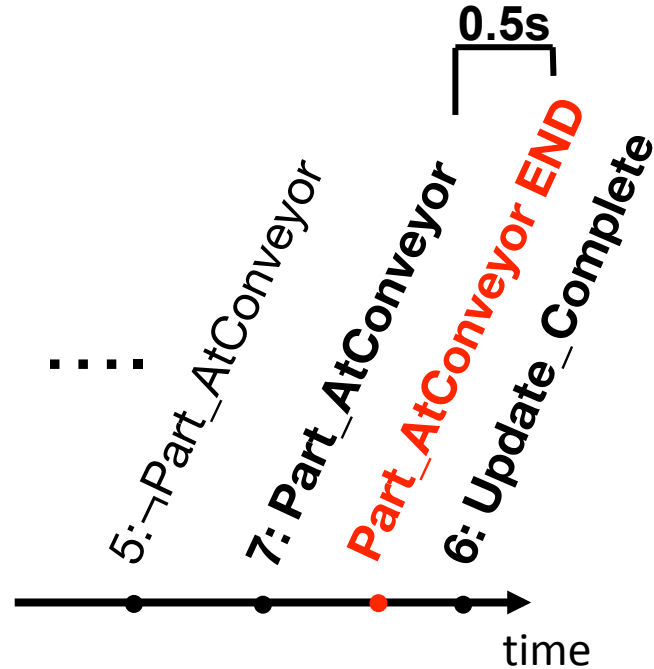
Violated TPTL Spec: $\square t_x.(\text{Pallet} \rightarrow \diamond t_y.(\text{Retract} \wedge t_y - t_x \leq 30s))$

Traditional Event Permutation **Doesn't** Solve the Problem

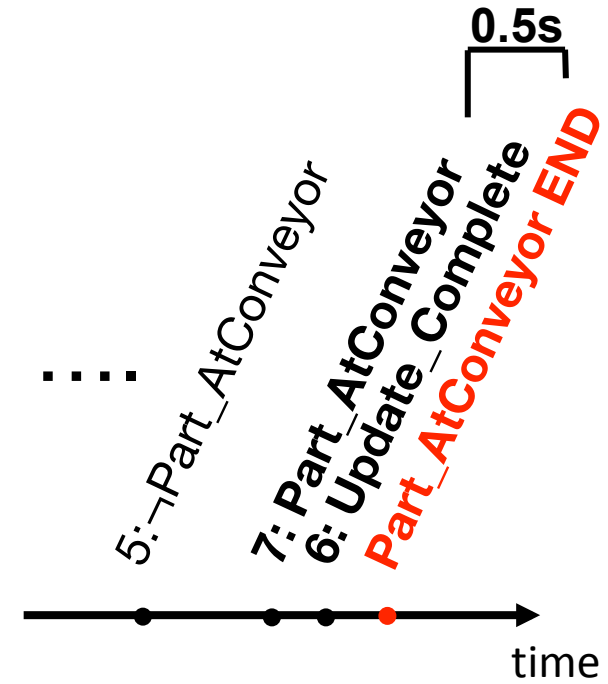
1->....->5->6->7 **Correct!**



5->7->6 **Error!**



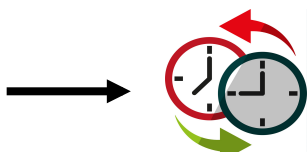
5->7->6 **Still Correct!**



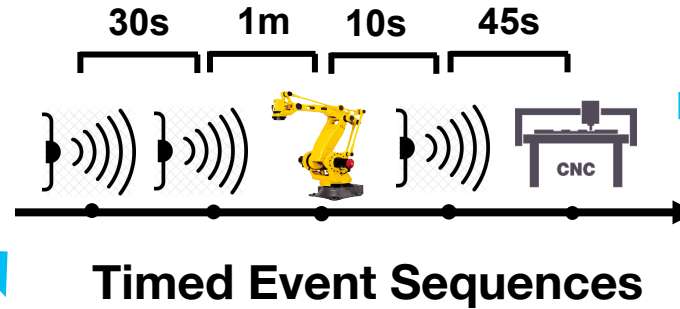
VETPLC: Generating **Timed Event Sequences** to enable Automated Safety Vetting of PLC Code



Program Analysis on PLC/Robot:
Generating Event Causality Graph



Data Mining on Runtime Data:
Discovering Temporal Invariants

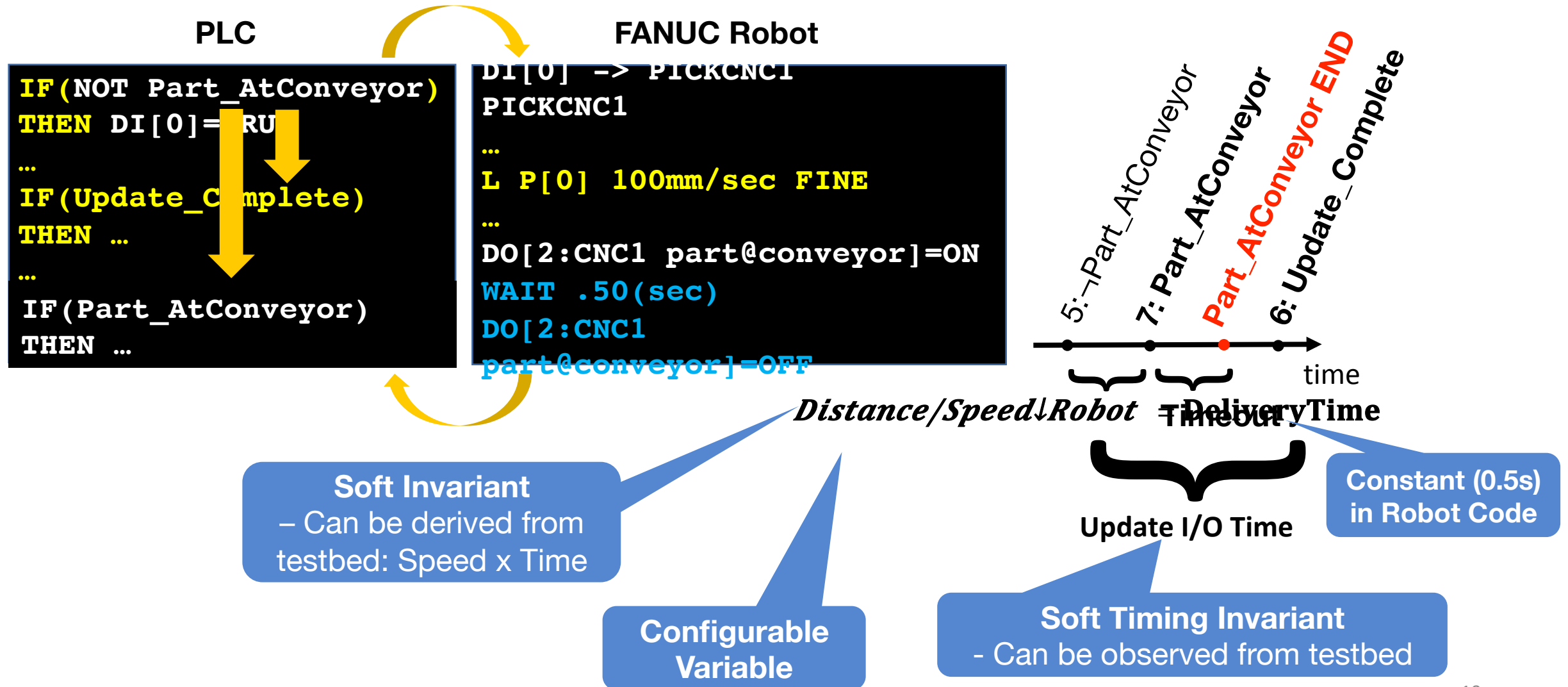


Execution Traces

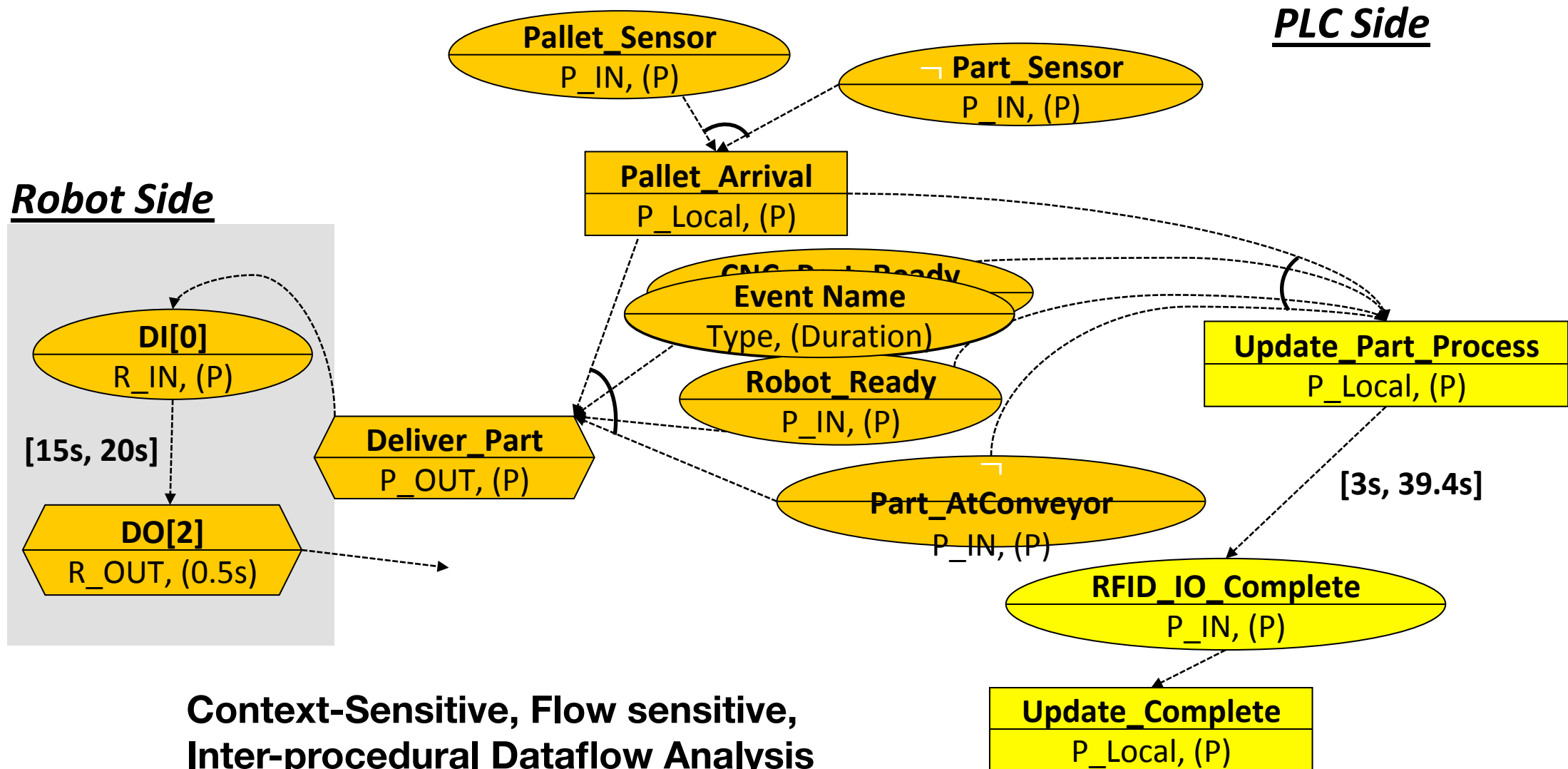


Safety Violations

VETPLC on Running Example



Timed Event Causality Graph (TECG): Find Valid Event Orders



Mining Temporal Invariants for Events: 2 Steps

Step 1: Qualitative “followed-by”:
– Synoptic (*FSE’11*)

Follows $[\epsilon_a][\epsilon_b] = \text{Occurrence}[\epsilon_a]$

Step 2: Quantitative “with-in”:
– Perfume (*ASE’14*)

$\square t_x.(\epsilon_a \rightarrow \Diamond t_y.(\epsilon_b \wedge t_y - t_x \geq \tau_{\text{lower}}))$

$\square t_x.(\epsilon_a \rightarrow \Diamond t_y.(\epsilon_b \wedge t_y - t_x \leq \tau_{\text{upper}}))$

Results for Motivating Example
(1.2 GB data for 10 hours):

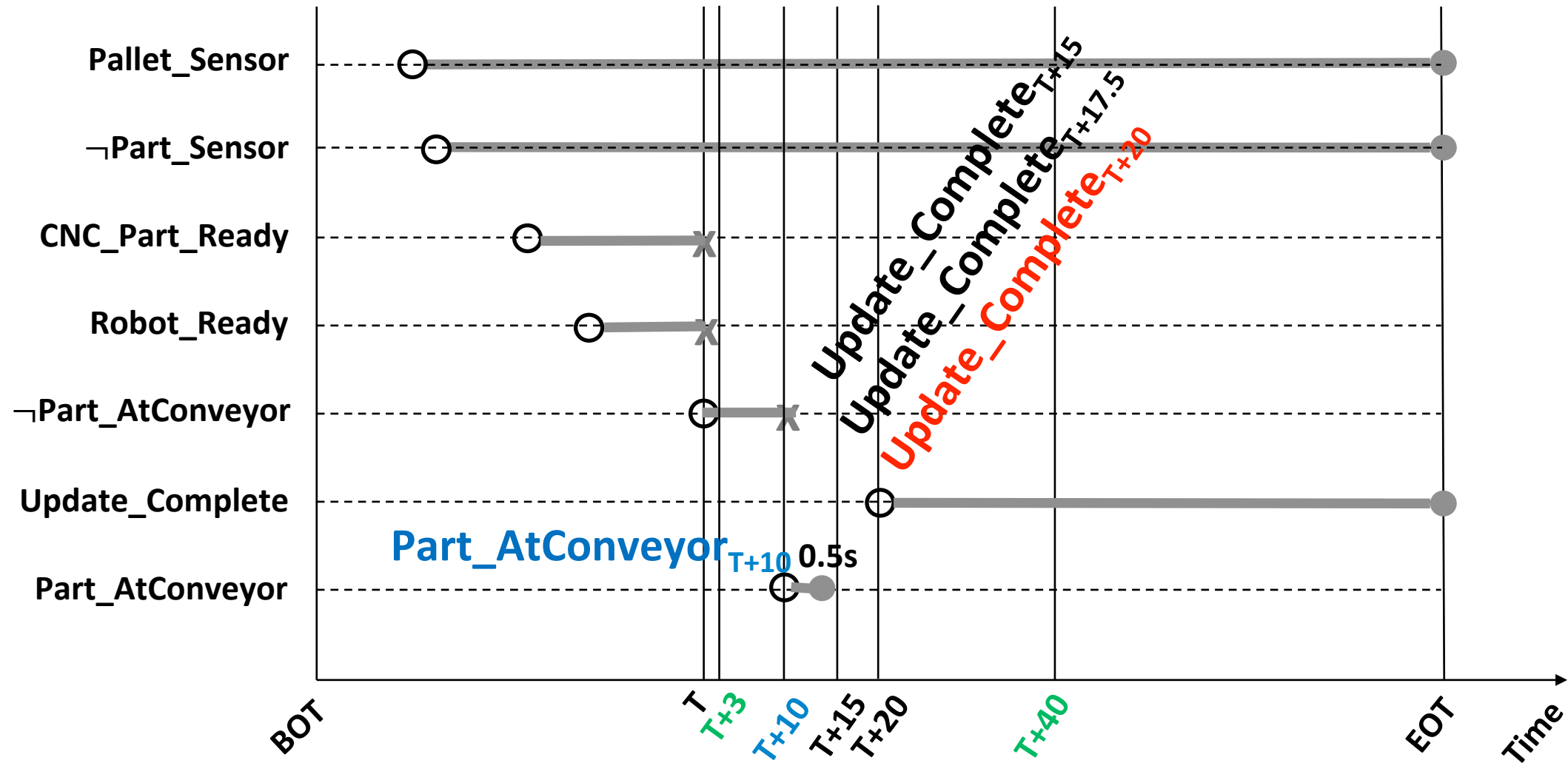
TABLE I: Mined Invariants

Event Pair	Invariant
$\square(\text{Deliver_Part} \rightarrow \Diamond \text{Part_AtConveyor})$	[24.4s, 24.6s]
$\square(\text{Update_Part_Process} \rightarrow \Diamond \text{RFID_IO_Complete})$	[15s, 20s]
$\square(\text{Update_Part_Process} \rightarrow \Diamond \text{Update_Complete})$	[15s, 20s]

Advantage of TECG: Only need to mine relations that do not contradict TECG

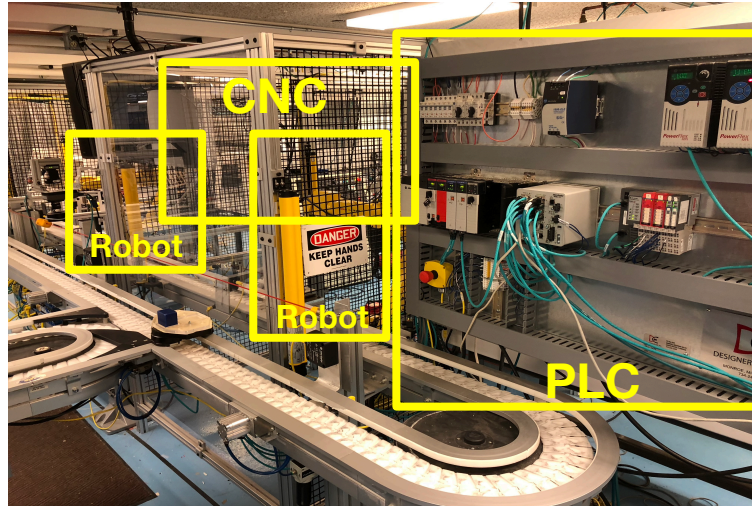
Creating Timed Event Sequences

Safety Violation Triggered
How to discretize durations?



Evaluation on Real Testbeds for Different Scenarios

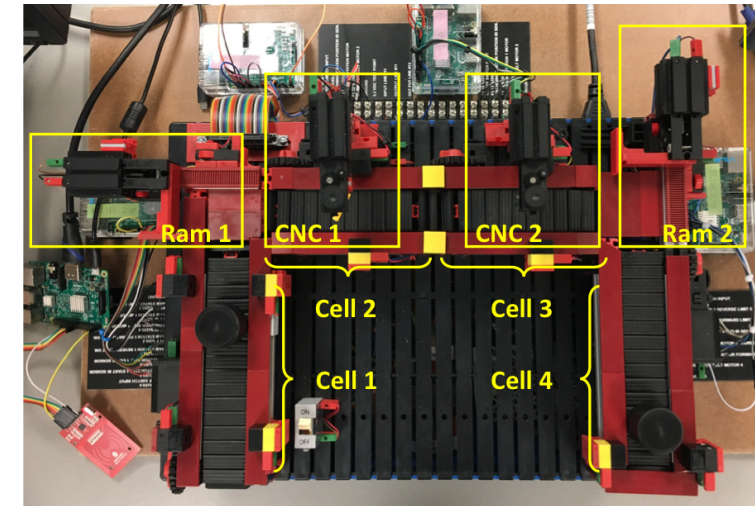
2 Different Testbeds



SMART: Automotive Production Line

10 Safety-critical Scenarios

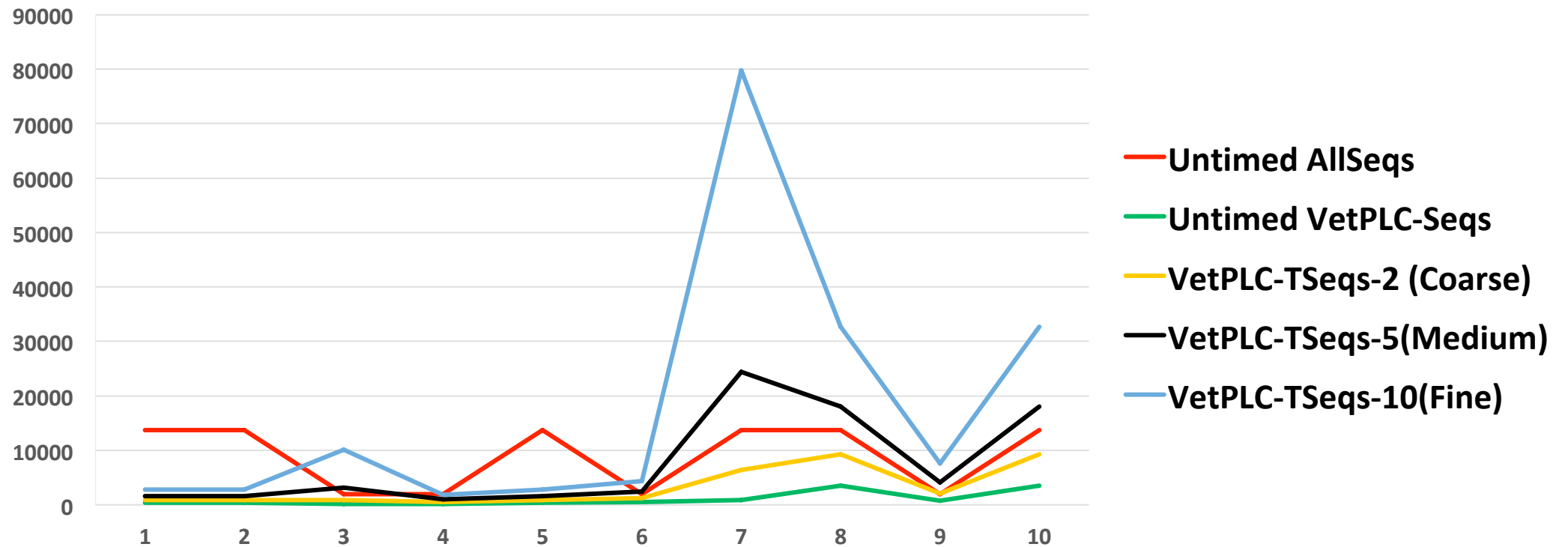
- S1: Conveyor Overflow #1
- S2: Robot in Danger Zone
- S3: Conveyor Overflow #2
- S4: Part-Gate Collision
- S5: CNC Overflow



Fischertechnik: Part Processing w/ 4 PLCs

- S6: Ram-Part Collision
- S7: CNC-Part Collision
- S8: Conveyor Overflow #3
- S9: Conveyor Underflow
- S10: Ram-Part Collision #2

Evaluation: How many sequences are created?



Red → Green: Program analysis reduces amount of event sequences

Green → Orange → Black → Blue: Time discretization can significantly increases that

Bug Detected? State-of-the-Art vs. VETPLC

State-of-the-art

VETPLC

#
1
2
3
4
5
6
7
8
9
10



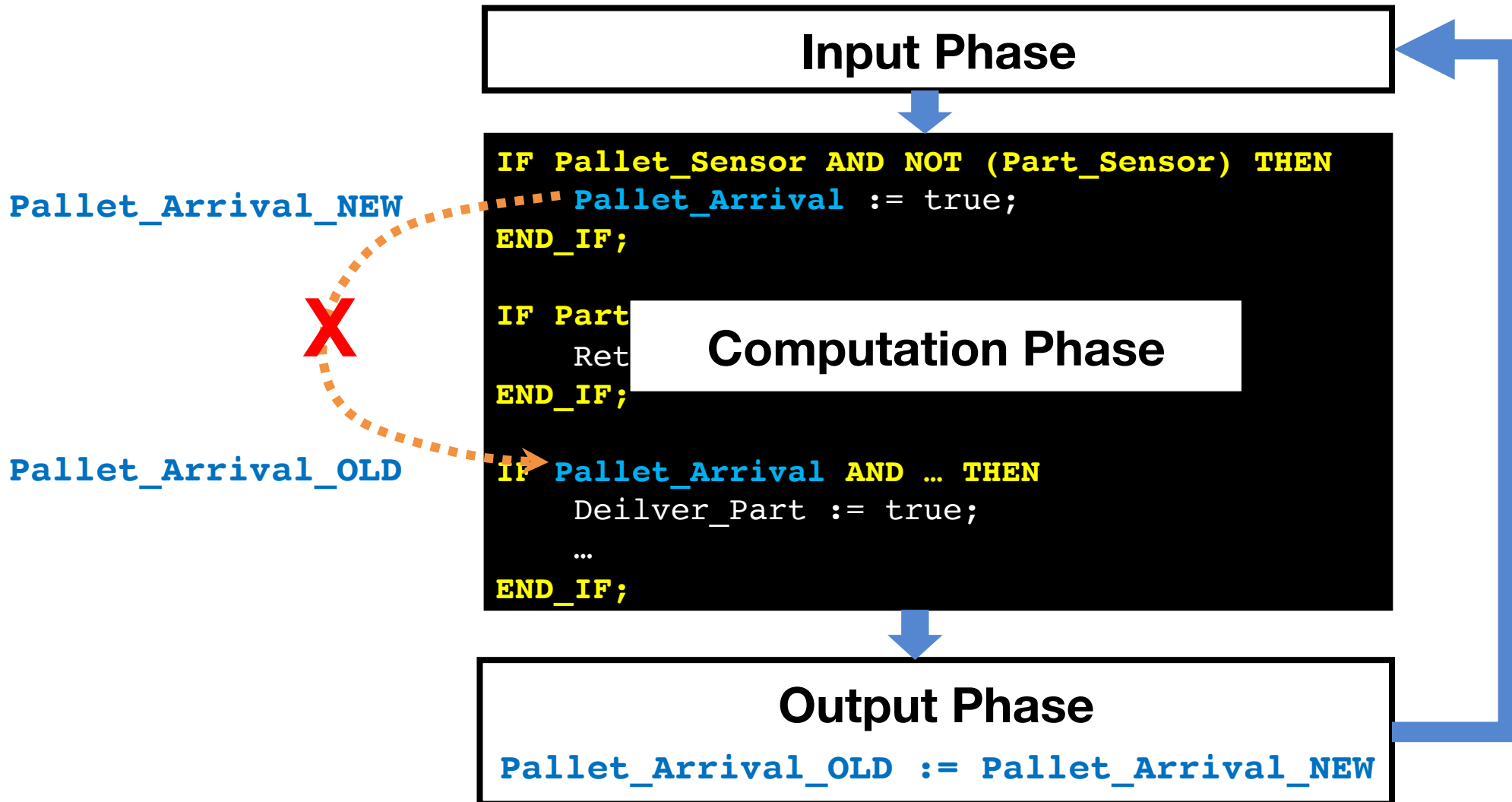
More Time for State-Error-Triggering Range
VETPLC Outputs More Precise State-Error-Triggering Range
Empirically, 5 slices works better.

Conclusion

- ❑ **Insight:** real-world PLC code is *event-driven* and *timing-sensitive*
- ❑ **Solution:** V_{ET}PLC *automatically constructs timed event sequences* via analyzing event causalities in PLC/robot code plus mining runtime data from physical testbeds
- ❑ **Effectiveness:** V_{ET}PLC outperforms state-of-the-art and has found “organic” vulnerabilities in two different types of real-world ICS testbeds.

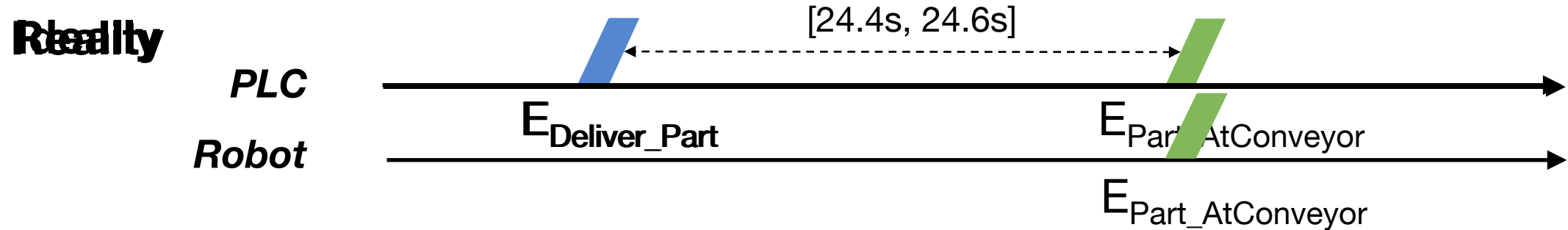
Thank you!

PLC Programming Paradigm: Scan Cycle

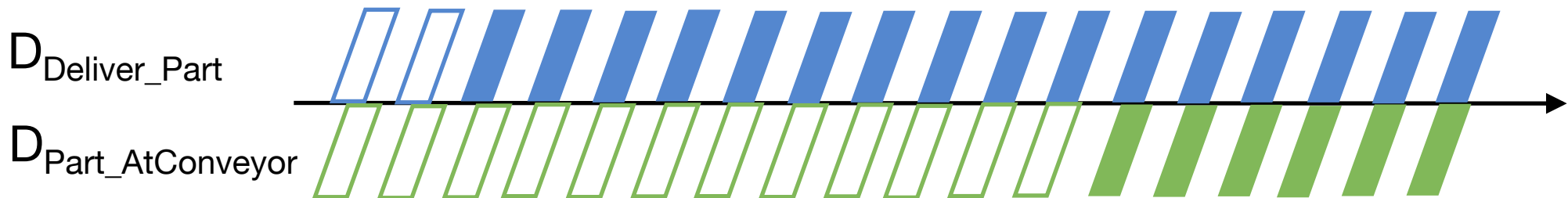


- ❖ No dataflow in one cycle
- ❖ Dataflow across cycles
- ❖ Any “Define” in a cycle may affect “Use” in the next

Technical Challenge: Distributed Event Sources



Solution: Inferring Events from State Variables



Speed Reconfiguration

∴

$\tau_{\downarrow lower} \leq T_{\downarrow job} = job / speed_{\downarrow conf} \leq \tau_{\downarrow upper}$ Time variation caused by **physical operations or program execution paths**

$speed_{\downarrow min} \leq speed_{\downarrow conf} \leq speed_{\downarrow max}$ Time variation caused by **reconfiguring machine speeds**

∴

$\tau_{\downarrow lower} \times speed_{\downarrow conf} / speed_{\downarrow max} \leq T_{\downarrow job} \leq \tau_{\downarrow upper} \times speed_{\downarrow conf} / speed_{\downarrow min}$

Speed_{rated}

Speed_{high-throughput-and-safe}



resulting in unpredictable robot motion and velocities.

3.41 slow speed control: A mode of **robot** motion control where the speed is limited to ≤ 250 mm/sec (10 in/sec) to allow persons sufficient time to either withdraw from **hazardous motion** or stop the robot.

3.42 space: The three dimensional volume encompassing the movements of all **robot** parts

Scenario-Specific Safety Specs

#	Description of Hazard	Safety Specification to Avoid Hazard
1	Motivating Example. See Section III	$\Box t_x.(\text{Pallet} \rightarrow \Diamond t_y.(\text{Retract_Stopper} \wedge t_y - t_x \leq 30\text{s}))$
2	Robot fails to return its safe zone.	$\Box t_x.(\neg \text{Safe_Zone} \rightarrow \Diamond t_y.(\text{Safe_Zone} \wedge t_y - t_x \leq 60\text{s}))$
3	Robot stops processing parts from conveyor due to signal conflicts.	$\Box t_x.(\text{Pallet} \rightarrow \Diamond t_y.(\text{Retract_Stopper} \wedge t_y - t_x \leq 30\text{s}))$
4	A pallet collides with a closed gate.	$\Box(\text{Pallet_AtGate} \rightarrow \Box \text{Gate_Open})$
5	CNC stops processing parts from gantry due to missing signals.	$\Box t_x.(\text{Part_In} \rightarrow \Diamond t_y.(\text{Part_Out} \wedge t_y - t_x \leq 5\text{m}))$
6	A ram starts pushing when a part has not fully entered the ram.	$\Box(\text{Part_Entering} \rightarrow \neg \Diamond \text{Ram_Push})$
7	A part is passed to CNC when a preceding part is not fully discharged.	$\Box(\text{CNC_Busy} \rightarrow \neg \Diamond \text{Part_Arrival})$
8	Parts are pushed to conveyor prematurely.	$\Box t_x.(\text{Part_Arrival} \rightarrow \Diamond t_y.(\text{Part_Arrival} \wedge t_y - t_x \leq 6\text{s}))$
9	A conveyor belt halts operation.	$\Box t_x.(\text{Part_Arrival} \rightarrow \Diamond t_y.(\text{Part_Arrival} \wedge t_y - t_x \geq 8.5\text{s}))$
10	Ram1 pushes a part to unprepared Ram2.	$\Box(\text{Part_Entering} \rightarrow \Box \text{Ram_Ready})$